

Introduction to the z/Architecture Transactional-Execution Facility

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Introduction

Introduced on the IBM System z EC12, the Transactional-Execution Facility is a suite of hardware instructions which allow multiple updates to be made to storage which either appear as a single operation or do not appear to occur as observed by programs. A series of instructions can be grouped together into a transaction and any instructions which are used to update memory will be committed to memory at the transaction's end.

Why is the Transactional-Execution Facility Important?

The toughest challenge for programmers in a multi-processor system is the issue of concurrency. The programmer is completely unaware of the order of execution between two threads and this can only be determined whilst the threads are running by the processor. If two threads are operating on the same data or are waiting for events to occur, it is important to serialise the operations of the threads so that the program's logic is able to guarantee the state of the threads. This is normally performed by using software-based locking methods which involve the operating system. As a result, a program that uses software-based locking can be much slower than one that relies on hardware locks. This also means that the program is dependent on the services and APIs provided by the operating system, possibly reducing the portability of the program.

The z/Architecture also provides a number of hardware-locking instructions such as COMPARE AND SWAP (which can be used to attempt an atomic update to memory) and PERFORM LOCKED OPERATION (which is an incredibly complex instruction).

What is the Transactional-Execution Facility used for?

The Transactional-Execution Facility can be used to provide:

- Atomicity – Transactions are either committed or aborted. Committed transactions will store all their updates to memory at transaction end. Aborted transactions will have no updates to memory stored at transaction end (except for non-transactional stores)
- Performance – Using the Transactional-Execution Facility can be much quicker than using software-based locking methods. If a program is permitted to use the Transactional-Execution Facility, then no operating system interaction is needed in order to have “hardware-style locking” governing a series of updates to memory.
- Reducing software complexity – series of COMPARE AND SWAP-style instructions can be replaced by code encapsulated inside a transaction.
- Improved diagnostics – The TDB (Transaction Diagnostics Block) can be examined in order to determine why a transaction was aborted and the updates to memory failed.

What is provided by the Transactional-Execution Facility?

The Transactional-Execution Facility provides the following:

- Three special-purpose controls – the Transaction Abort PSW, the Transaction Diagnostic Block address and the Transaction Nesting Depth.
- Five control register bits – these are used generally by the control program to govern what happens during an interrupt, what happens when transactions are aborted and if the Transactional-Execution Facility is available for programs to use.
- Six general purpose instructions:
 - TBEGIN – Start a non-constrained transaction
 - TBEGINC – Start a constrained transaction
 - TABORT – Abort a transaction
 - TEND – End a transaction

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- NTSTG – Non-Transactional Store
- ETND – Extract Transaction Nesting Depth

Prerequisites for using the Transactional-Execution Facility

In order to use the EC12's transactional memory facility, the machine must be equipped with the Transactional-Execution facility installed in hardware. Facility indicator bit 73 indicates whether the Transactional-Execution Facility is available. If both bits 73 and 50 are on, then the Constrained Transactional-Execution Facility is installed.

In order to use the instructions provided by the Transactional-Execution Facility, the PTFs for APAR PM49761 (or RPM asma90-1.6.0-18.rpm or higher) must be applied to HLASM.

Assembling a program that uses the Transactional-Execution Facility

Assembly of a program that uses the Transactional-Execution Facility is the same as that for any program. However, HLASM provides the listing exit ASMAXTXP which can be used to check whether the program being assembled violates any of the constrained transaction rules. Please note that this is an *assembly time* check and not a *runtime* check and therefore is limited to checking code as if it was to be run sequentially.

ASMAXTXP will check code for:

- Use of a zero base register for constrained transactions
- Allow AR modification control being 0 for constrained transactions
- Whether or not a transaction exceeds its instruction and/or byte limit
- Restricted instructions between the issue of a TBEGINC and TEND instruction

To enable the ASMAXTXP exit, use HLASM with the option `EXIT (PRTEXTIT (ASMAXTXP))`.

If any of the above checks fail, then the transaction may not execute correctly due to transaction violations. ASMAXTXP will issue a warning message for any failing check with a return code 4.

Terminology

Constrained Transaction

A constrained transaction is started by the TBEGINC instruction. Updates to memory will be committed upon the end of the transaction specified by TEND. Constrained transactions are limited to 256 bytes in length between the TBEGINC and TEND instructions, many instructions are classified as *restricted* and are not permitted inside a constrained transaction, and no more than 32 instructions may be executed by the transaction.

Non-constrained Transaction

A non-constrained transaction is started by the TBEGIN instruction. At the end of the transaction specified by the TEND instruction, an attempt will be made to commit any updates to memory. If the updates to memory are unsuccessful, the transaction will be *aborted*; control passes to the instruction following the TBEGIN and the condition code is set to non-zero.

If a non-constrained transaction is aborted, general purpose registers specified by the GRSM (General-purpose Register Save Mask) operand of the TBEGIN instruction are restored to their previous values.

Non-transactional Store

The NTSTG instruction allows a store to memory to be performed whilst inside a transaction in a non-transactional way. This means that if the transaction is aborted, the non-transactional store would still have been committed.

Using the Transactional-Execution Facility

The rest of this Technote demonstrates how to use the Transactional-Execution Facility by changing a program which currently uses software-based locking methods to update a linked list.

The program consists of:

- A master process which performs the program's initialisation, spawns child processes and waits for them to finish their processing before terminating the program.
- Each child process creates a linked-list which are chained together with the lists from other child processes. In order to do this, the child processes currently use the z/OS ISGENQ macros to obtain and release software locks to perform safe updates to the chain of linked-lists.
- After attaching its linked-list into the chain, the child process waits for a predetermined amount of time in order to simulate processing after which it goes through the chain removing its linked-list.

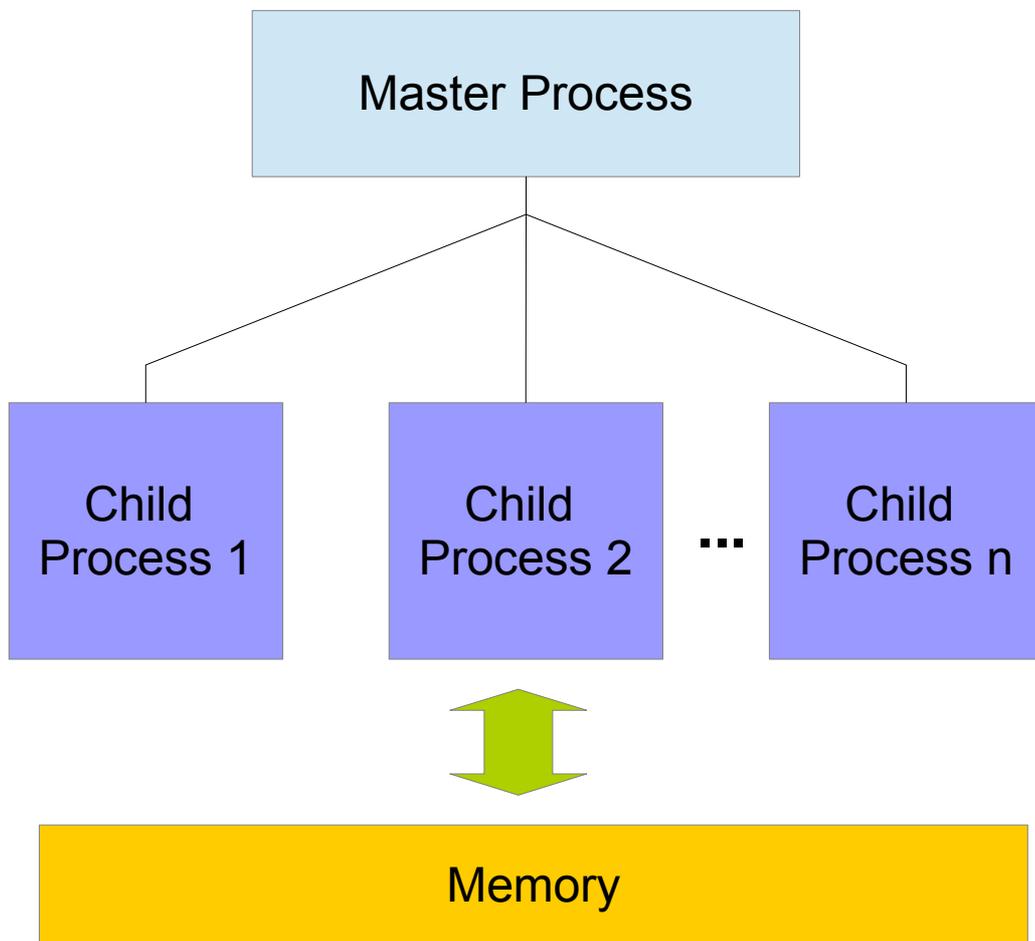


Figure 1 – Example master and child processes

Changing the master process

During initialisation the master process can check whether or not the Transactional-Execution Facility is available for use. This can be done by examining the output of the STFLE instruction. Under z/OS, the PSA already stores the output of the STFLE instruction.

```

-----1-----2-----3-----4-----5-----6-----7-----8
* INCLUDE THE IHAPSA MACRO IN ORDER TO EXAMINE THE PSA
  IHAPSA ,
* Address PSA and examine facility bits
  XR    R0,R0
  USING PSA,R0
  TM    FLCEFACILITIESLISTBYTE9,X'40' Test for trans facility
  IF    (O)                               Is it on?
      MVI  TB_TRAN_SUPPORT,C'Y'           Yes - non-trans enabled
      TM    FLCEFACILITIESLISTBYTE6,X'20' Constr trans available?
      IF    (O)
          MVI  TB_TRAN_SUPPORT,C'C'       Yes - constr trans on
      ENDIF ,
  ELSE ,
      MVI  TB_TRAN_SUPPORT,C'N'           Indicate no trans
  ENDIF ,
  DROP  R0

* Each child process can now examine the data passed to it in
* TB_TRAN_SUPPORT to determine whether or not it should run in
* transactional, constrained or non-transactional mode. It is
* important that all child processes run in the same mode so that no
* problems occur should a child process using the Transactional-Execution
* Facility does not get into a conflict with a child process using
* software based locking.

```

At this stage of the program, each child process will be able to examine the state of TB_TRAN_SUPPORT to determine which “mode” to run in. The program has been written for portability so that the same object code can run on machines with and without the Transactional-Execution Facility being available. The value of TB_TRAN_SUPPORT is one of:

- C → Constrained transactions are available for use
- Y → Transaction-Execution facility is available for use
- N → No transaction facility is available for use

Child process code – Chaining the lists together

Each child process creates its own linked-list in memory. This part of the processing does not affect other child processes since these lists are created independently. Once the list is created, it must be chained together with the other linked-lists in memory and it is this operation which must be conducted in an atomic way to prevent other child processes attempting to update the chain as they add their list at the same time.

```

-----+-----1-----+-----2-----+-----3-----+-----4-----+-----5-----+-----6-----+-----7-----+-----8
* Examine WS_TRAN_SUPPORT (child working storage copy of
* TB_TRAN_SUPPORT) and start the chaining operation.
  CLI  WS_TRAN_SUPPORT,C'C'
  IF   (EQ)
    J   CTRACT          Run the constrained transaction
  ENDIF ,
  CLI  WS_TRAN_SUPPORT,C'N'
  IF   (EQ)
    XR   R5,R5          Clear transaction retry count
    LARL R4,NTRACTS     Address of transaction retry
    J    NTRACT         Run non-constrained transaction
  ELSE ,
    GET  ENQ             Get ENQ and software lock
  ISQENQ REQUEST=OBTAIN,
                                X
        QNAME=S_ENQ_MAJOR_NAME,
                                X
        RNAME=S_ENQ_MINOR_NAME,
                                X
        RNAMELEN=S_ENQ_RNAME_LENGTH,
                                X
        CONTROL=EXCLUSIVE,
                                X
        ENQTOKEN=WS_ENQ_TOKEN,
                                X
        SCOPE=STEP,
                                X
        MF=(E,WS_ISGENQ_S)
    J   COM_Q           Skip transaction initialisation
  ENDIF ,

```

The child process can now either initialise its transaction or begin the common chaining code. The common chaining code is at the label COM_Q.

Starting the non-constrained transaction

```

-----+-----1-----+-----2-----+-----3-----+-----4-----+-----5-----+-----6-----+-----7-----+-----8
* Registers 6,7,8 and 9 will be modified during the transaction.
* Should the transaction abort for any reason, these should be
* restored to their previous values.
REGS69 EQU B'00011000'      Create the GRSM for regs 6-9
*
* Start the non-constrained transaction
*
NTRACTS TBEGIN WS_TRAN_TDB,REGS69
        JNZ  TRAN_ABORT     Setup abort handler
        J    COM_Q         Jump to common code

```

Starting a non-constrained transaction involves setting up the transaction abort routine, providing a Transaction Diagnostic Block (TDB) and also the General-purpose Register Save Mask (GRSM). If the transaction is aborted, control will pass to the instruction immediately after the TBEGIN instruction. This instruction is used to jump to the common transaction abort routine TRAN_ABORT. The general-purpose registers indicated in the GRSM will be restored to their values previous to their values prior to the transaction starting.

Starting the constrained transaction

```

-----1-----2-----3-----4-----5-----6-----7-----8
*
* Start a constrained transaction
*
CTRACT    TBEGINC 0,0
          JNZ      TRAN_ABORT    Setup abort handler if we don't start
    
```

If the constrained transaction facility is available to the program, then the program will make use of it since using a constrained transaction *always* commits any updates it makes. However, constrained transactions have far more limitations placed upon them both in terms of which instructions may be used within them and the number of instructions which are to be placed inside them.

Common chaining code

Once the transactions have been started, or the ENQ has been obtained for the critical section of code, then the common chaining code can begin:

```

-----1-----2-----3-----4-----5-----6-----7-----8
*
* Common chaining code for chaining 2 lists into 2 chains
*
COM_Q    L        R8,TB_CELL_CHAIN_1    Get head 1 into working storage
          L        R9,TB_CELL_CHAIN_2    Get head 2 into working storage
          ST       R8,WS_CELL_1HEAD
          ST       R9,WS_CELL_2HEAD
          ST       R6,TB_CELL_CHAIN_1    Save new head 1
          ST       R7,TB_CELL_CHAIN_2    Save new head 2
          L        R6,WS_CELL_1TAIL      Get last cell 1
          L        R7,WS_CELL_2TAIL      Get last cell 2
          ST       R8,SP1_F_PTR          Join pointer
          ST       R9,SP2_F_PTR
          L        R8,TB_ENQ_COUNT_1     Get chain counts
          L        R9,TB_DEQ_COUNT_1
          AFI      R8,1                   Increase chain counts
          AFI      R9,1
          ST       R8,TB_ENQ_COUNT_1
          ST       R9,TB_DEQ_COUNT_1
          L        R8,TB_CELL_COUNT_1
          L        R9,TB_CELL_COUNT_2
          A        R8,WS_CELL_COUNT_1
          A        R9,WS_CELL_COUNT_2
          ST       R8,TB_CELL_COUNT_1
          ST       R9,TB_CELL_COUNT_2
*
* Check if running a transaction and if so then end it.
* If not then release the ENQ
*
          CLI      WS_TRAN_SUPPORT,C'N'
          IF       (NE)
              TEND                        End the transaction to commit
          ELSE    ,                        Release the ENQ
          ISGENQ  REQUEST=RELEASE,        X
                  ENQTOKEN=WS_END_TOKEN, X
                  MF=(E,WS_ISGENQ_S)
          ENDIF    ,
    
```

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In the common code, a number of counters and pointers were updated inside the transaction. Doing the same operation without a software-based lock and without the transactional-execution facility would become very difficult.

If changing existing code to make use of transactional-execution rather than using other software-based locking methods, critical section code may need to be restructured so that it can be executed within a transaction. The difficulty in doing this should be considered by the programmer before attempting to modify their existing code.

More information about the restrictions of using the Transactional-execution Facility can be found in the z/Architecture Principles of Operations (SA22-7832-09) – Chapter 5 – Program Execution.

Creating a Transaction-Abort Routine

If a transaction fails to initialise or it is aborted either by the hardware or by issuing the TABORT instruction, control passes to the instruction immediately following the corresponding TBEGIN / TBEGINC instruction and the condition code is set to non-zero. This is why in the example program, a JNZ TRAN_ABORT is used immediately following the TBEGIN and TBEGINC instructions.

The transaction abort routine may include logic for examining the Transaction Diagnostic Block (TDB). The format of the TDB is documented in Principles of Operation and contains many useful fields for diagnosing why a transaction was aborted and the state of the Transaction-Execution Facility at the time of the abort such as the Transaction Nesting Depth (TND), Transaction Abort Code (TAC), Aborted-Transaction Instruction Address (AITA).

The PERFORM PROCESSOR ASSIST (PPA) instruction can be used during the transaction abort processing in order to try and increase the likelihood of the transaction committing successfully. Before starting the non-constrained transaction, the child process uses register 5 as a transaction retry count and register 4 as the transaction retry address.

The transaction abort routine code is as follows:

```

-----+-----1-----+-----2-----+-----3-----+-----4-----+-----5-----+-----6-----+-----7-----+-----8
* *****
* Transaction abort routine TRAN_ABORT
* *****
*
* If the transaction fails to initialise or if it is aborted
* then control will pass here.
* The TDB may contain information that will help to diagnose
* why the transaction failed.
* The number of times that the transaction has failed is
* stored in register 5.
* The transaction retry address is stored in register 4.
*
TRAN_ABORT DC      0H
             WTO      MF=(E,WS_WTO_TABORT)  Show transaction has aborted
             CLI      WS_TRAN_SUPPORT,C'Y'
             IF      (NE)
*
             SOMETHING WENT SERIOUSLY WRONG!!
             LA       R3,WS_TRAN_TDB        Make TDB easy to access
             ABEND    20,DUMP                ABEND and dump
             ENDIF   ,
*
* All the code for examining the TDB for any useful information
* should go here...
*
* Update the transaction fail count
             AHI      R5,1
             CIJNL   R5,20,DISTRAN          Failed 20 times? - Yes, give up
             PPA     R5,0,1                  No - ask for help
             BR      R4                      Redrive transaction
DISTRAN MVI     WS_TRAN_SUPPORT,C'N'        Disable transaction support
             J       TIDYUP                  Go tidy up storage and exit

```

When to use transactions

Although the Transaction-Execution Facility provides a very useful, they cannot be used to solve every locking problem. Transactions are good for short updates to memory that require more than a single instruction. Transactions should not be used for preparing to enter a critical section, timing / synchronising threads and large updates to memory.

Transactions and software-based locking methods should not be used together for the same critical sections of code. Since the Transaction-Execution Facility is based in hardware, it takes no notice of any software-based locks that are managed by the operating system. Consider the following:

- Processes A and B are attempting to update the same set of counters in memory.
- Process A acquires a software-based lock for the region of memory with exclusive access.
- Process B does not wait for the lock and instead uses a transaction to update the memory and commits its changes.
- Process A (which has the software lock) writes its changes to memory after Process B had finished its transaction and therefore overwrites the changes that Process B had made corrupting the data.

The problem occurs because Process B is unaware to the software-based lock held by Process A and therefore acts as if no locking was used over the data.

Summary

The Transactional-Execution Facility provided by System z can be used to simplify code and improve performance in the correct circumstances when needing to provide a level of atomicity to small critical sections of code.

Constrained transactions have severe limitations placed upon them, they always commit when they reach the TEND instruction and so can be used to guarantee the state of data if the transaction started.

Non-constrained transactions do not suffer with the same limitations as constrained transactions however they may be aborted if they cannot commit their data and therefore require a suitable transaction abort routine to be prepared should this occur. Any general-purpose registers specified in the GRSM (General-purpose Register Save Mask) upon entry to the non-constrained transaction will be restored to their previous values. The transaction abort routine can use the PERFORM PROCESSOR ASSIST (PPA) instruction in order to increase the likelihood of a non-constrained transaction committing its changes to data.

Appendix – JCL and Program Source Code

This technote provides a downloadable zip file which contains the source code, JCL and the assembler options file XOPT.txt

Contents of the zip file

\JCL	
\JCL\RUN.txt	← JCL used to run the executable load module
\JCL\TRNSASX.txt	← JCL used to build the executable load module
\SOURCE	
\SOURCE\TBCHILD.asm	← Child program source
\SOURCE\TBDATA.asm	← Common data copybook
\SOURCE\TBMSTR.asm	← Master program source
\XOPT.txt	← HLASM options data set

Required data sets on z/OS

The following data sets should be created on your z/OS system under a suitable high-level qualifier before uploading the contents of the zip file to the system:

&myhlq..HLASM.TRNSACTN.JCL

 This is the PDS into which the JCL is to be uploaded

&myhlq..HLASM.TRNSACTN.LOAD

 This PDS will contain the executable load module built by the JCL

&myhlq..HLASM.TRNSACTN.SOURCE

 This is the PDS into which the source files are to be uploaded

&myhlq..HLASM.TRNSACTN.XOPT

 This sequential data set will contain the HLASM options.

Contents of HLASM options data set

The HLASM options data set (specified by the ASMAOPT DD statement) allows for a convenient place for assembly options to be located.

```
INFO
OBJ
EXIT (PRTEXTIT (ASMAXTXP) )
```

In particular interest for building programs that use the transactional-execution facility, the HLASM-supplied print exit ASMAXTXP should be used to check that your program adheres to the restrictions of the transactional-execution facility.

For more information on using ASMAXTXP, refer to the Appendix P in the HLASM Programmer's Guide (SC26-4941-06).

JCL used to assemble programs

The program is designed to assemble and issue WTOs so that it is clear to see how the program operates. To disable this, use an alternative library which has an inert WTO macro in it during assembly.

Note that the sample program makes use of the HLASM Structured Programming Macros which are available as part of the High Level Assembler Toolkit Feature. See the High Level Assembler Release 6 Toolkit Feature Users' Guide (GC26-8710-10) for more information at: <http://www.ibm.com/software/awdtools/hlasm/library.html>

```

-----+-----1-----+-----2-----+-----3-----+-----4-----+-----5-----+-----6-----+-----7-----+-----8
//TRNSASX JOB NOTIFY=&SYSUID
//*
//* *****
//* THIS JCL IS USED TO BUILD AND LKED AN ASSEMBLER PROGRAM.
//* THE FOLLOWING IS ASSUMED WHEN USING THIS JCL:
//* THE SOURCE FOR THE PROGRAM RESIDES IN ----> &SRCLIB
//* THE EXECUTABLE LOADMOD IS OUTPUT TO -----> &LODLIB
//* THE LISTING IS OUTPUT TO -----> &PRFXNM..Z.&PRGNM
//* THE PROGRAM NAME IS -----> &PRGNM
//* THE ASSEMBLER OPTIONS USED ARE IN -----> &PRFXNM..XOPT
//* *****
//*
//* CHANGE THE FOLLOWING LINE TO REFLECT THE PROGRAM NAME
// SET PRGNM=TBMSTR
//*
//* CHANGE THE FOLLOWING LINE TO REFLECT THE HLQ FOR THE SOURCE
// SET PRFXNM=&SYSUID..HLASM.TRNSACTN
//*
// SET SRCLIB=&PRFXNM..SOURCE
// SET LODLIB=&PRFXNM..LOAD
// SET OPTLIB=&PRFXNM..XOPT
//*
//* CHANGE THE FOLLOWING TO REFLECT THE HLQ FOR HLASM AND SYSTEM
//* DATA SETS
// SET HLQ=SYS1
// SET HLQASM=SYS1.ASM
//*
//* COMPILE AND LINK STEP - PRODUCE THE EXECUTABLE LOAD MOD
//COMPLNK PROC
//ASMSAMP EXEC ASMACL
//STEPLIB DD DISP=SHR,DSN=&HLQASM..SASMOD1
// DD DISP=SHR,DSN=&HLQASM..SASMOD2
//SYSLIB DD DISP=SHR,DSN=&SRCLIB
// DD DISP=SHR,DSN=&HLQ..MACLIB
// DD DISP=SHR,DSN=&HLQASM..SASMMAC1
// DD DISP=SHR,DSN=&HLQASM..SASMMAC2
//C.SYSIN DD DISP=SHR,DSN=&SRCLIB (&PRGNM)
//C.SYSPRINT DD DISP=OLD,DSN=&PRFXNM..Z.&PRGNM
//C.ASMASOPT DD DISP=SHR,DSN=&OPTLIB
//L.SYSLMOD DD DISP=SHR,DSN=&LODLIB (&PRGNM)
// PEND
//* *****
// SET PRGNM=TBMSTR
//CP1 EXEC COMPLNK
//* *****
// SET PRGNM=TBCHILD
//CP2 EXEC COMPLNK
//* *****

```

TBDATA – DSECT used by TBMSTR and TBCHILD

The TBDATA DSECT is used by both the master (TBMSTR) and child (TBCHILD) programs for providing a common data structure in which the master program can communicate with each child program created.

```

-----+-----1-----+-----2-----+-----3-----+-----4-----+-----5-----+-----6-----+-----7-----+-----8
* *****
* TBDATA - DSECTs for TBMMASTER and TBCHILD programs
*
* Copyright IBM(UK) Ltd 2014
* HLASM - High Level Assembler and Toolkit (5696-234)
*
* IBM HLASM - Transactional-Execution Facility Technote
*
* This COPYBOOK is used to provide the TBMSTR and TBCHILD programs
* with a common set of data structures.
*
* The TB_MASTER DSECT is the data area which is passed from the
* master program to the child programs.
*
* The format of memory areas obtained from the 1st and 2nd subpool is
* defined by the TB_SP1_PARM and TB_SP2_PARM DSECTs and are used by
* the child programs in order to build storage chains.
*
* *****
TB_MASTER      DSECT ,           Area passed to children tasks
TB_EYE         DC CL(C_EYE_LENGTH)' ' Eye catcher
TB_CELL_ID_1   DC F'0'           cell poolid
TB_CELL_ID_2   DC F'0'           cell poolid
TB_ENQ_COUNT_1 DC F'0'           Number of ...
TB_DEQ_COUNT_1 DC F'0'           Number of ...
TB_CELL_COUNT_1 DC F'0'          Number of cells
TB_CELL_COUNT_2 DC F'0'          Number of cells
TB_ATTA_COUNT  DC F'0'           Attach count
TB_DETA_COUNT  DC F'0'           Dettach count
TB_CELL_CHAIN_1 DC A(0)          cell poolid
TB_CELL_CHAIN_2 DC A(0)          cell poolid
TB_TRAN_SUPPORT DC C'N'          Trans support availability
              DC 0D              align eye
TB_FCLTY_LIST  DC 2FD'0'         Faciliy list
TB_EYE_E       DC CL(C_EYE_LENGTH)' ' Eye catcher
              DC 0D
TB_LENGTH      EQU *-TB_MASTER
* *****
TB_SP1_PARM    DSECT ,           Subpool 1
SP1_EYE        DC CL(C_EYE_LENGTH)' ' Eye catcher
SP1_TCB        DC A(0)           Pointer to TCB
SP1_F_PTR      DC A(0)           Pointer to next cell
              DC 0D
SP1_LENGTH     EQU *-TB_SP1_PARM
* *****
TB_SP2_PARM    DSECT ,           Subpool 2
SP2_EYE        DC CL(C_EYE_LENGTH)' ' Eye catcher
SP2_TCB        DC A(0)           Pointer to TCB
SP2_F_PTR      DC A(0)           Pointer to next cell
              DC 0D
SP2_LENGTH     EQU *-TB_SP2_PARM
    
```

TBMSTR – The master program

The purpose of the master program is to detect the presence of the transactional-execution facility and create child processes which create and maintain memory subpools.

The value of C_CHILD_COUNT determines the number of child processes to create.

By uncommenting the lines in the master program marked with the comment ##### will cause the master program not to determine the presence of the transactional-execution facility.

```

-----+-----1-----+-----2-----+-----3-----+-----4-----+-----5-----+-----6-----+-----7-----+-----8
TBMSTR   RSECT
TBMSTR   RMODE   ANY
TBMSTR   AMODE   31
* #####
* TBMSTR - Master transaction program
*
* Copyright IBM(UK) Ltd 2014
* HLASM - High Level Assembler and Toolkit (5696-234)
*
* The purpose of this program is to spawn child processes in order to
* build linked lists.
* The number of child processes spawned is determined by the
* C_CHILD_COUNT equate.
* #####
*       COPY ASMMSP
*       ASMDREG
*       ASMMREL
C_EYE_LENGTH EQU 16
C_CHILD_COUNT EQU 100           Amount of child threads to create
*       PRINT OFF
*       SYSSTATE ARCHLVL=1
*       IEABRCX DEFINE
*       PRINT ON
*       IHAPSA ,
*
* #####
* Start of mainline code
* #####
*
TBMSTR   RSECT ,
*       PRINT ON
*       SAVE (14,12), , *
*       LARL R9,STATICAREA
*       USING (STATICAREA,STATICAREA_E),R9
*       LA R2,WS_LENGTH
*       STORAGE OBTAIN,LENGTH=(R2),ADDR=(R3),COND=NO
*       USING (WORKING_STORAGE,WORKING_STORAGE+WS_LENGTH),R3
*       ST R13,WS_SAVEA+4      save callers savearea
*       LA R13,WS_SAVEA       address my savearea
*       ST R13,WS_SAVEA+8     .. and save it
*       MVC WS_EYE,=CL(C_EYE_LENGTH)'TBMSTR W/S->'
*       MVC WS_EYE_E,=CL(C_EYE_LENGTH)'TBMSTR W/S-<'
*       USING TB_MASTER,WS_MASTER_DATA
*       MVC TB_EYE,=CL(C_EYE_LENGTH)'TBMSTR MASTER>'
*       MVC TB_EYE_E,=CL(C_EYE_LENGTH)'<TBMSTR MASTER'
*       MVI TB_TRAN_SUPPORT,C'N'
*       XC TB_FCLTY_LIST,TB_FCLTY_LIST
*       MVC WS_LOAD(S_LOAD_L),S_LOAD

```

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```

-----1-----2-----3-----4-----5-----6-----7-----8
*
* Setup WTOs showing transactional execution facility
*
      MVC   WS_WTO_SUPPORT_Y(S_WTO_SUPPORT_YL),S_WTO_SUPPORT_Y
      MVC   WS_WTO_SUPPORT_C(S_WTO_SUPPORT_CL),S_WTO_SUPPORT_C
      MVC   WS_WTO_SUPPORT_N(S_WTO_SUPPORT_NL),S_WTO_SUPPORT_N
      MVC   WS_WTO_ENDED(S_WTO_ENDED_L),S_WTO_ENDED
*
      MVC   WS_ATTACHX(S_ATTACHX_L),S_ATTACHX
      MVC   WS_POOL_1(S_POOL_1_L),S_POOL_1
      MVC   WS_POOL_2(S_POOL_2_L),S_POOL_2
*
* Check for the availability of the transaction facility
* This is located in bits 73 and 50 of the result of the STFLE
* instruction.
* Bit 73 must be checked first and indicates the presence of the
* Transactional-Execution Facility.
* Bit 50 may be checked once it is know that the transactional
* execution facility is available on this machine and indicates the
* Constrained transactional execution facility is available.
*
* If you wish to build the program so that it does not use the
* transactional execution facility then uncomment the line marked
* #####
*      AGO   .B_LOAD                               #####
      XR    R0,R0
      USING PSA,R0
      TM    FLCEFACILITIESLISTBYTE9,X'40'         Test for trans facility
      IF    (O)                                     Is it on?
          MVI TB_TRAN_SUPPORT,C'Y'                 Yes - non-constr enabl'd
          WTO MF=(E,WS_WTO_SUPPORT_Y)
          TM    FLCEFACILITIESLISTBYTE6,X'20'     Constr trans available?
          IF    (O)
              MVI TB_TRAN_SUPPORT,C'C'           Yes - constr trans on
              WTO MF=(E,WS_WTO_SUPPORT_C)
          ENDIF ,
      ELSE ,
          CLI TB_TRAN_SUPPORT,C'N'                 Unknown value?
          JNE CABEND                               Yes - ABEND
          WTO MF=(E,WS_WTO_SUPPORT_N)             Indicate no tran
      ENDIF ,
      DROP R0
.B_LOAD ANOP ,                                     #####
*
* Load and start the child programs TBCHILD.
* After the child programs have been started, TBMSTR will wait until
* the child programs finish their processing.
*
BEGLOAD  LOAD  EP=TBCHILD,                          X
          SF=(E,WS_LOAD)
          ST   R0,WS_TBCHILD_EP
          ST   R1,WS_TBCHILD_LP

```

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```

-----+-----1-----+-----2-----+-----3-----+-----4-----+-----5-----+-----6-----+-----7-----+-----8
CPOOL BUILD, PCELLCT=10, SCELLCT=5,                                X
      CSIZE=128, SP=0, LOC=31,                                    X
      CPID=TB_CELL_ID_1,                                        X
      HDR='TBMSTR CPOOL ONE',                                  X
      MF=(E, WS_POOL_1)
CPOOL BUILD, PCELLCT=8, SCELLCT=12,                               X
      CSIZE=64, SP=0, LOC=31,                                  X
      CPID=TB_CELL_ID_2,                                        X
      HDR='TBMSTR CPOOL TWO',                                  X
      MF=(E, WS_POOL_2)
XR     R5, R5
LA     R6, C_CHILD_COUNT
LA     R7, WS_ECB_CHILD
LA     R8, WS_TCB_CHILD
NextAttach DC 0H
      ATTACHX EP=TBCHILD, PARAM=TB_MASTER,                       X
              ECB=(R7),                                         X
              SF=(E, WS_ATTACHX), MF=(E, WS_TBCHILD_PL)
      ST     R1, 0(, R8)
      LA     R7, L'WS_ECB_CHILD(, R7)
      LA     R8, L'WS_TCB_CHILD(, R8)
      AFI   R5, 1
      BCT   R6, NextAttach
      ST     R5, TB_ATTA_COUNT
      XR     R5, R5
      LA     R6, C_CHILD_COUNT
      LA     R7, WS_ECB_CHILD
      LA     R8, WS_TCB_CHILD
NextEndAndDetach DC 0H
      WAIT  1, ECB=(R7)
      DETACH (R8)
      LA     R7, L'WS_ECB_CHILD(, R7)
      LA     R8, L'WS_TCB_CHILD(, R8)
      AFI   R5, 1
      BCT   R6, NextEndAndDetach
      ST     R5, TB_DETA_COUNT
CLEANUP DC 0H
      DELETE EPLOC=WS_TBCHILD_EP
      CPOOL DELETE,                                             X
              CPID=TB_CELL_ID_1
      CPOOL DELETE,                                             X
              CPID=TB_CELL_ID_2
*
* At this stage, the child processes have completed processing.
* TBMSTR releases its storage and returns control to the OS.
*
EXIT   DC      0H                      exit label
      WTO     MF=(E, WS_WTO_ENDED)
      L       R13, WS_SAVEA+4          restore callers savearea
      LA     R2, WS_LENGTH
      STORAGE RELEASE, LENGTH=(R2), ADDR=(R3), COND=NO
      RETURN (14, 12), , RC=0         return
CABEND DC 20X'0'

```

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```

-----1-----2-----3-----4-----5-----6-----7-----8
*
* *****
* Static storage used by TBMSTR
* *****
*
STATICAREA          DC 0D'0'                align storage
*
* Static macro definitions
*
S_LOAD              LOAD  SF=L
S_LOAD_L            EQU   *-S_LOAD
S_ATTACHX           ATTACHX SF=L
S_ATTACHX_L         EQU   *-S_ATTACHX
S_POOL_1            CPOOL BUILD,MF=L
S_POOL_1_L          EQU   *-S_POOL_1
S_POOL_2            CPOOL BUILD,MF=L
S_POOL_2_L          EQU   *-S_POOL_2
*
S_CHILD_NAME        DC   CL8'TBCHILD'
*
* WTO messages for static storage
*
S_WTO_SUPPORT_Y     WTO  'Transactionality available',MF=L
S_WTO_SUPPORT_YL    EQU  *-S_WTO_SUPPORT_Y
S_WTO_SUPPORT_C     WTO  'Constrained Transactionality available',MF=L
S_WTO_SUPPORT_CL    EQU  *-S_WTO_SUPPORT_C
S_WTO_SUPPORT_N     WTO  'No transactionality available',MF=L
S_WTO_SUPPORT_NL    EQU  *-S_WTO_SUPPORT_C
S_WTO_ENDED         WTO  'Master process ended',MF=L
S_WTO_ENDED_L       EQU  *-S_WTO_ENDED
*
                LTORG ,
STATICAREA_E        DC 0D'0'                align storage - end of static
*
* *****
* Working storage used by TBMSTR
* *****
*
WORKING_STORAGE     DSECT                    Program dynamic area
WS_EYE              DC CL(C_EYE_LENGTH)' '   Eye catcher
WS_SAVEA            DC 18F'00'               save area
WS_LOAD             DC 0D,XL(S_LOAD_L)'00'   load macro
WS_ATTACHX          DC 0D,XL(S_ATTACHX_L)'00' Attachx
WS_POOL_1           DC 0D,XL(S_POOL_1_L)'00' cpool 1
WS_POOL_2           DC 0D,XL(S_POOL_2_L)'00' cpool 2
*
* WS WTO buffers for transaction support
*
WS_WTO_SUPPORT_Y    DC 0D,XL(S_WTO_SUPPORT_YL)'00'
WS_WTO_SUPPORT_C    DC 0D,XL(S_WTO_SUPPORT_CL)'00'
WS_WTO_SUPPORT_N    DC 0D,XL(S_WTO_SUPPORT_NL)'00'
WS_WTO_ENDED        DC 0D,XL(S_WTO_ENDED_L)'00'
*
WS_TBCHILD_EP       DC A(0)                  child program address
WS_TBCHILD_LP       DC A(0)                  child program address

```

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```

-----1-----2-----3-----4-----5-----6-----7-----8
*
* *****
* Static storage used by TBMSTR
* *****
*
STATICAREA          DC 0D'0'                align storage
*
* Static macro definitions
*
S_LOAD              LOAD  SF=L
S_LOAD_L            EQU   *-S_LOAD
S_ATTACHX          ATTACHX SF=L
S_ATTACHX_L        EQU   *-S_ATTACHX
S_POOL_1           CPOOL BUILD,MF=L
S_POOL_1_L         EQU   *-S_POOL_1
S_POOL_2           CPOOL BUILD,MF=L
S_POOL_2_L         EQU   *-S_POOL_2
*
S_CHILD_NAME       DC   CL8'TBCHILD'
*
* WTO messages for static storage
*
S_WTO_SUPPORT_Y    WTO  'Transactionality available',MF=L
S_WTO_SUPPORT_YL   EQU  *-S_WTO_SUPPORT_Y
S_WTO_SUPPORT_C    WTO  'Constrained Transactionality available',MF=L
S_WTO_SUPPORT_CL   EQU  *-S_WTO_SUPPORT_C
S_WTO_SUPPORT_N    WTO  'No transactionality available',MF=L
S_WTO_SUPPORT_NL   EQU  *-S_WTO_SUPPORT_C
S_WTO_ENDED        WTO  'Master process ended',MF=L
S_WTO_ENDED_L     EQU  *-S_WTO_ENDED
*
                LTORG ,
STATICAREA_E       DC 0D'0'                align storage - end of static
*
* *****
* Working storage used by TBMSTR
* *****
*
WORKING_STORAGE    DSECT                    Program dynamic area
WS_EYE             DC CL(C_EYE_LENGTH)' '   Eye catcher
WS_SAVEA           DC 18F'00'              save area
WS_LOAD            DC 0D,XL(S_LOAD_L)'00'   load macro
WS_ATTACHX         DC 0D,XL(S_ATTACHX_L)'00' Attachx
WS_POOL_1          DC 0D,XL(S_POOL_1_L)'00' cpool 1
WS_POOL_2          DC 0D,XL(S_POOL_2_L)'00' cpool 2
*
* WS WTO buffers for transaction support
*
WS_WTO_SUPPORT_Y   DC 0D,XL(S_WTO_SUPPORT_YL)'00'
WS_WTO_SUPPORT_C   DC 0D,XL(S_WTO_SUPPORT_CL)'00'
WS_WTO_SUPPORT_N   DC 0D,XL(S_WTO_SUPPORT_NL)'00'
WS_WTO_ENDED       DC 0D,XL(S_WTO_ENDED_L)'00'
*
WS_TBCHILD_EP     DC A(0)                  child program address
WS_TBCHILD_LP     DC A(0)                  child program address

```

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```
-----1-----2-----3-----4-----5-----6-----7-----8
WS_TBCHILD_PL  DC A(0)                child program address
WS_TCB_CHILD   DC (C_CHILD_COUNT)F'00'    TCB
WS_ECB_CHILD   DC (C_CHILD_COUNT)F'00'    ECB
WS_COUNTER     DC F'00'                ws counter
               DC 0D                    align end
WS_MASTER_DATA DC XL(TB_LENGTH)'00'
WS_EYE_E       DC CL(C_EYE_LENGTH)' '    Eye catcher
WS_LENGTH      EQU *-WORKING_STORAGE WS Length
               COPY TBDATA
               END   TBMSTR
```

TBCHILD – The child program

The purpose of the child program is to build linked-lists of acquired storage. Each child process is created by the master program and has passed to it a number of parameters. As the child begins to start its processing, it uses the WS_TRAN_SUPPORT field in order to determine whether it is able to make use of the transactional-execution facility.

```

-----+-----1-----+-----2-----+-----3-----+-----4-----+-----5-----+-----6-----+-----7-----+-----8
TBCHILD  RSECT
TBCHILD  RMODE  ANY
TBCHILD  AMODE  31
          COPY ASMMSP
          ASMDREG
          ASMMREL
* *****
* TBCHILD - Child transaction program
*
* Copyright IBM(UK) Ltd 2014
* HLASM - High Level Assembler and Toolkit (5696-234)
*
* The purpose of this program is to link a number of storage areas
* into a shared linked list.
*
* *****
C_EYE_LENGTH  EQU 16
C_POOL1_COUNT EQU 16          Amount of cells to create in pool1
C_POOL2_COUNT EQU 32          Amount of cells to create in pool2
C_LOOP_COUNT  EQU 1024
          PRINT OFF
          SYSSTATE ARCHLVL=1
          IEABRCX  DEFINE
          PRINT ON
*
* Prepare program execution environment
*
          SAVE (14,12),,*
          LARL  R9,STATICAREA
          USING (STATICAREA,STATICAREA_E),R9
          L    R10,0(,R1)      Point to input parms
          USING TB_MASTER,R10
*
* Ensure that the data we have been passed by the caller is correct
*
          CLC   TB_EYE,=CL(C_EYE_LENGTH) 'TBMSTR MASTER>'
          BNE  E_IS_NOK
          CLC   TB_EYE_E,=CL(C_EYE_LENGTH) '<TBMSTR MASTER'
          BE   E_IS_OK
E_IS_NOK DC   0H
          DC   A(9)
E_IS_OK  DC   0H
          LA   R2,WS_LENGTH
          STORAGE OBTAIN,LENGTH=(R2),ADDR=(R3),COND=NO
          USING (WORKING_STORAGE,WORKING_STORAGE+WS_LENGTH),R3
          ST   R13,WS_SAVEA+4   save callers savearea
          LA   R13,WS_SAVEA     address my savearea
          ST   R13,WS_SAVEA+8   .. and save it

```

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```

-----+-----1-----+-----2-----+-----3-----+-----4-----+-----5-----+-----6-----+-----7-----+-----8
*
* Prepare working storage eyecatchers
*
      MVC   WS_EYE,=CL(C_EYE_LENGTH)'TBCHILD W/S->'
      MVC   WS_EYE_E,=CL(C_EYE_LENGTH)'TBCHILD W/S-<'
      MVC   WS_ISGENQ_S(S_ISGENQ_S_L),S_ISGENQ_S
      MVC   WS_WTO_ENDED(S_WTO_ENDED_L),S_WTO_ENDED
*
* Setup transaction support messages
*
      MVC   WS_TRAN_SUPPORT,TB_TRAN_SUPPORT
      MVC   WS_WTO_SUPPORT_Y(S_WTO_SUPPORT_YL),S_WTO_SUPPORT_Y
      MVC   WS_WTO_SUPPORT_C(S_WTO_SUPPORT_CL),S_WTO_SUPPORT_C
      MVC   WS_WTO_SUPPORT_N(S_WTO_SUPPORT_NL),S_WTO_SUPPORT_N
      MVC   WS_WTO_TABORT(S_WTO_TABORT_L),S_WTO_TABORT
      MVC   WS_WTO_DIS_TRAN(S_WTO_DIS_TRAN_L),S_WTO_DIS_TRAN
      MVC   WS_WTO_TEND(S_WTO_TEND_L),S_WTO_TEND
      MVC   WS_WTO_TNCS(S_WTO_TNCS_L),S_WTO_TNCS
      MVC   WS_WTO_TNCE(S_WTO_TNCE_L),S_WTO_TNCE
      MVC   WS_WTO_TCS(S_WTO_TCS_L),S_WTO_TCS
      MVC   WS_WTO_TCE(S_WTO_TCE_L),S_WTO_TCE
      MVC   WS_WTO_WAITS(S_WTO_WAITS_L),S_WTO_WAITS
      MVC   WS_WTO_WAITE(S_WTO_WAITE_L),S_WTO_WAITE
      MVC   WS_TRAN_EYE_S,=CL(C_EYE_LENGTH)'TRANSACTION>>>>'
      MVC   WS_TRAN_EYE_E,=CL(C_EYE_LENGTH)'<<<<TRANSACTION'
      XC    WS_TRAN_TDB(C_TDB_LEN),WS_TRAN_TDB    Clear the TDB
*
      USING PSA,0
      MVC   WS_PSATOLD,PSATOLD
      DROP  0
*
* Start of mainline code
*
      LA    R11,C_LOOP_COUNT
StartProcess DC 0H
      ST    R11,WS_LOOP_CNTR          Save loop counter
      XR    R5,R5
      LA    R7,WS_CELL_1P_GET
      LA    R8,C_POOL1_COUNT
*
* Start building a chain from the 1st subpool
*
GetPool1 DC 0H
      USING TB_SP1_PARM,R1
      CPOOL GET,                      X
      CPID=TB_CELL_ID_1
      ST    R1,0(,R7)
      LA    R7,4(,R7)
      AFI   R5,1
      MVC   SP1_EYE,=CL(C_EYE_LENGTH)'TBCHILD SP1->'
      MVC   SP1_TCB,WS_PSATOLD
      BCT   R8,GetPool1
      DROP  R1
      ST    R5,WS_CELL_COUNT_1
      XR    R5,R5
      LA    R7,WS_CELL_2P_GET
      LA    R8,C_POOL2_COUNT

```

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```

-----1-----2-----3-----4-----5-----6-----7-----8
*
* Start building a chain from the 2nd subpool
*
GetPool2 DC    0H
         USING TB_SP2_PARM,R1
         CPOOL GET,                                     X
         CPID=TB_CELL_ID_2
         ST    R1,0(,R7)
         LA    R7,4(,R7)
         AFI   R5,1
         MVC   SP2_EYE,=CL(C_EYE_LENGTH)'TBCILD SP2->'
         BCT   R8,GetPool2
         DROP  R1
         ST    R5,WS_CELL_COUNT_2
*
* Chain the cells together
*
         LA    R6,WS_CELL_1P_GET
         LA    R8,C_POOL1_COUNT
         BCTR  R8,0
         USING TB_SP1_PARM,R7
Chain1  DC    0H
         L     R7,0(,R6)          current cell
         LA    R6,4(,R6)          get next one
         ST    R6,SP1_F_PTR       save forward cell
         BCT   R8,Chain1         next cell
         L     R7,0(,R6)          last cell
         ST    R7,WS_CELL_1TAIL  save tail cell
         DROP  R7
         LA    R6,WS_CELL_2P_GET
         LA    R8,C_POOL2_COUNT-1
         USING TB_SP2_PARM,R7
Chain2  DC    0H
         L     R7,0(,R6)          current cell
         LA    R6,4(,R6)          get next one
         ST    R6,SP2_F_PTR       save forward cell
         BCT   R8,Chain2         next cell
         L     R7,0(,R6)          last cell
         ST    R7,WS_CELL_2TAIL  save tail cell
         DROP  R7
         LA    R6,WS_CELL_1P_GET
         LA    R7,WS_CELL_2P_GET
         USING TB_SP1_PARM,R6
         USING TB_SP2_PARM,R7
         L     R6,0(,R6)          get new head
         L     R7,0(,R7)          get new head
*
* Check for the presence of the transaction facility
* If we can run a constrained transaction on the data rather
* than using a full ENQ, then do so.
*
         LR    R2,R9              Free up register 9
         DROP  R9
         USING (STATICAREA,STATICAREA_E),2
         CLI   WS_TRAN_SUPPORT,C'C'

```

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```

-----1-----2-----3-----4-----5-----6-----7-----8
      IF      (EQ)                                Run constr transaction?
      WTO     MF=(E,WS_WTO_SUPPORT_C)
      J       CTRACT
      ENDIF ,
      CLI     WS_TRAN_SUPPORT,C'Y'              Run non-constr transaction?
      IF      (EQ)
      WTO     MF=(E,WS_WTO_SUPPORT_Y)
      J       NTRACT
      ENDIF ,
* If control flows through here, then run non-transactional code
GETENQ     WTO     MF=(E,WS_WTO_SUPPORT_N)
          ISGENQ  REQUEST=OBTAIN,                X
          QNAME=S_ENQ_MAJOR_NAME,              X
          RNAME=S_ENQ_MINOR_NAME,              X
          RNAMELEN=S_ENQ_RNAME_LENGTH,         X
          CONTROL=EXCLUSIVE,                    X
          ENQ_TOKEN=WS_ENQ_TOKEN,              X
          SCOPE=STEP,                            X
          MF=(E,WS_ISGENQ_S)
          J       COM_Q                          Skip transaction initialisation
*
* The transactional code will modify registers 6,7,8 and 9.
* Should the transaction abort for any reason, then these registers
* will be restored to their previous values.
REGS69     EQU    b'00011000'                  Transaction GRSM for registers 6-9
*
* Start a non-constrained transaction
*
NTRACT     XR      R5,R5                        Set register 5 as trns abort count
          LARL    R4,NTRACTS                    Set address of transaction retry
          LARL    R11,GETENQ                    Set address of non trans retry
NTRACTS    TBEGIN WS_TRAN_TDB,REGS69
          JNZ     TRAN_ABORT                    Setup abort handler
          J       COM_Q                          Jump to common code
*
* Start a constrained transaction
*
CTRACT     TBEGINC 0,0
          JNZ     TRAN_ABORT                    Setup abort handler
COM_Q      L      R8,TB_CELL_CHAIN_1            Get head 1 into working storage
          L      R9,TB_CELL_CHAIN_2            Get head 2 into working storage
          ST     R8,WS_CELL_1HEAD
          ST     R9,WS_CELL_2HEAD
          ST     R6,TB_CELL_CHAIN_1            save head
          ST     R7,TB_CELL_CHAIN_1            save head
          L      R6,WS_CELL_1TAIL              get last cell
          L      R7,WS_CELL_2TAIL              get last cell
          ST     R8,SP1_F_PTR                    join chain
          ST     R9,SP2_F_PTR                    join chain
          L      R8,TB_ENQ_COUNT_1
          L      R9,TB_DEQ_COUNT_1
          AFI    R8,1
          AFI    R9,1
          ST     R8,TB_ENQ_COUNT_1
          ST     R9,TB_DEQ_COUNT_1

```

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```

-----1-----2-----3-----4-----5-----6-----7-----8
    AFI    R8,1
    AFI    R9,1
    ST     R8,TB_ENQ_COUNT_1
    ST     R9,TB_DEQ_COUNT_1
    L      R8,TB_CELL_COUNT_1
    L      R9,TB_CELL_COUNT_2
    A      R8,WS_CELL_COUNT_1
    A      R9,WS_CELL_COUNT_2
    ST     R8,TB_CELL_COUNT_1
    ST     R9,TB_CELL_COUNT_2
*
* Check if we're running in a transaction and if we are not, then
* jump over the TEND instruction
*
    CLI    WS_TRAN_SUPPORT,C'N'
    IF     (NE)
* End the transaction and don't release an ENQ
    TEND
    WTO    MF=(E,WS_WTO_TEND)
    ELSE  ,
    DROP  R6
    DROP  R7
*
* Non-transactional ENQ release
*
    ISGENQ REQUEST=RELEASE,           X
           ENQTOKEN=WS_ENQ_TOKEN,    X
           MF=(E,WS_ISGENQ_S)
    ENDIF ,
*
* Simulate processing by waiting a short interval before continuing
*
    WTO    MF=(E,WS_WTO_WAITS)
MCONT    STIMER WAIT,TUINTVL=S_WAIT_TIME
    WTO    MF=(E,WS_WTO_WAITE)
*
* After the simulated processing wait is over, start releasing the
* storage and decoupling the chains.
*
*
* Prepare registers before iterating through
*
    MVC    WS_ISGENQ_S(S_ISGENQ_S_L),S_ISGENQ_S
    LA     R6,TB_CELL_CHAIN_1        Get top of chain
    LA     R8,C_POOL1_COUNT          number of entries to release
    LA     R7,WS_CELL_1P_REL         save pointers here
*
* Check the state of the transaction-execution facility and change
* into non-constrained transational mode if it is available.
*
    CLI    WS_TRAN_SUPPORT,C'N'
    IF     (NE)
        MVI    WS_TRAN_SUPPORT,C'Y'
    ENDIF ,
    CLI    WS_TRAN_SUPPORT,C'C'

```

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```

-----1-----2-----3-----4-----5-----6-----7-----8
      IF      (EQ)
        WTO      MF=(E,WS_WTO_TCS)
        J        CTRACTDQ
      ENDIF ,
      CLI      WS_TRAN_SUPPORT,C'Y'
      IF      (EQ)
        WTO      MF=(E,WS_WTO_TNCS)
        J        NTRACTDQ
      ENDIF ,
* Else run in non-transactional mode
DQENQ   ISGENQ REQUEST=OBTAIN,                X
        QNAME=S_ENQ_MAJOR_NAME,             X
        RNAME=S_ENQ_MINOR_NAME,             X
        RNAMELEN=S_ENQ_RNAME_LENGTH,        X
        CONTROL=EXCLUSIVE,                  X
        ENQTOKEN=WS_ENQ_TOKEN,              X
        SCOPE=STEP,                          X
        MF=(E,WS_ISGENQ_S)
        J        COM_DEQ
*
* Start dequeuing chain 1 items
*
      USING TB_SP1_PARM,R6
COM_DEQ DC      0H
      CLI      WS_TRAN_SUPPORT,C'N'   Are we running in non-trans mode?
      IF      (NE)
        CLI      WS_TRAN_SUPPORT,C'C'   Run in constr trans mode?
        IF      (NE)
NTRACTDQ XR      R5,R5                 No - clear trans count
          LARL   R4,NTRACTSQ           Prepare retry address
          LARL   11,DQENQ
NTRACTSQ TBEGIN WS_TRAN_TDB,REGS69     Start unconstrained trans
          JNZ   TRAN_ABORT             Prepare abort handler
        ELSE ,
CTRACTDQ TBEGINC 0,0
          JNZ   TRAN_ABORT
        ENDIF ,
      ENDIF ,
DQCOM   ST      R6,0(,R7)
        LA     R7,4(,R7)
        L      R6,SP1_F_PTR
        CLI      WS_TRAN_SUPPORT,C'N'
        IF      (NE)
          TEND
        ENDIF ,
        BCT    R8,COM_DEQ
        ST     R6,TB_CELL_CHAIN_1
        DROP   R6
        CLI      WS_TRAN_SUPPORT,C'N'
        IF      (NE)
          TEND
        ENDIF ,

```

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-----1-----2-----3-----4-----5-----6-----7-----8
*
* Start dequeuing chain 2 items
*
      L      R6,TB_CELL_CHAIN_2      get top of chain
      USING TB_SP2_PARM,R6
      LA     R8,C_POOL2_COUNT        number of entries to release
      LA     R7,WS_CELL_2P_REL save pointers here
UncPool2 DC     0H
      CLI   WS_TRAN_SUPPORT,C'N'     Non-transactional mode?
      IF    (NE)
          CLI   WS_TRAN_SUPPORT,C'C'   Constrained transactional mode?
          IF    (EQ)
              TBEGINC 0,0              Yes - start constrained trans
              JNZ   TRAN_ABORT
          ELSE ,
              TBEGIN WS_TRAN_TDB,REGS69 No - start unconstrained trans
              JNZ   TRAN_ABORT        Prepare abort handler
          ENDIF ,
      ENDIF ,
uncp2_c ST     R6,0(,R7)
      LA     R7,4(,R7)
      L      R6,SP2_F_PTR
      CLI   WS_TRAN_SUPPORT,C'N'
      IF    (NE)
          TEND
          BCT  R8,UncPool2             Branch and reinitialise trans
      ELSE ,
          BCT  R8,uncp2_c             Just loop back under the lock
      ENDIF ,
      ST     R6,TB_CELL_CHAIN_2
      DROP  R6
*
* If no transaction had taken place, then release the ENQ.
*
      CLI   WS_TRAN_SUPPORT,C'N'
      IF    (EQ)
RELDQENQ ISGENQ REQUEST=RELEASE,          X
          ENQTOKEN=WS_ENQ_TOKEN,          X
          MF=(E,WS_ISGENQ_S)
      ENDIF ,
*
* Now that the chains have been uncoupled from the main chain, the
* storage for them can be released.
*
MCONT2  LR     R9,R2                  Restore register 9 for Wrk Strg
      DROP  R2
      USING (STATICAREA,STATICAREA_E),R9
      LA     R7,WS_CELL_1P_GET
      LA     R8,C_POOL1_COUNT
FrePool1 DC     0H
      L      R6,0(,R7)
      CPOOL FREE,                      X
          CPID=TB_CELL_ID_1,           X
          CELL=(R6)
      LA     R7,4(,R7)

```

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-----1-----2-----3-----4-----5-----6-----7-----8
      BCT   R8,FrePool1
      LA    R7,WS_CELL_2P_GET
      LA    R8,C_POOL1_COUNT
FrePool2 DC    0H
      L     R6,0(,R7)
      CPOOL FREE,                                X
          CPID=TB_CELL_ID_2,                      X
          CELL=(R6)
      LA    R7,4(,R7)
      BCT   R8,FrePool2
*
* If the process has more work to do, then loop around and start again
*
      L     R11,WS_LOOP_CNTR      reload loop counter
      BCT   R11,StartProcess     loop around
*
* *****
* Mainline return point
* *****
EXIT   DC    0H                      exit label
      WTO   MF=(E,WS_WTO_ENDED)
      L     R13,WS_SAVEA+4        restore callers savearea
      LA    R2,WS_LENGTH
      STORAGE RELEASE,LENGTH=(R2),ADDR=(R3),COND=NO
      RETURN (14,12),,RC=0       return
*
* *****
* TRANSACTION ABORT ROUTINE
* *****
*
* If the transaction fails to initialise or if it is aborted,
* then control will pass here to the TRAN_ABORT label.
*
* If the program was running in a non-constrained transaction, then
* we examine the Transaction Diagnostic Block (TDB)
*
* The number of times that a transaction attempts to redrive itself
* is determined by register 5.
*
* The transactional retry address is stored in register 4.
*
* The non-transactional retry address (once register 5 has reached a
* certain threshold), is stored in register 11.
*
TRAN_ABORT DS  0H
          WTO   MF=(E,WS_WTO_TABORT)
          CLI   WS_TRAN_SUPPORT,C'Y'      Non-constrained transaction?
          IF    (NE)
*          SOMETHING WENT SERIOUSLY WRONG!!
          LA    R3,WS_TRAN_TDB
          ABEND 20,DUMP
DISTRAN MVI   WS_TRAN_SUPPORT,C'N'      Disable tran support for child
          WTO   MF=(E,WS_WTO_DIS_TRAN)
          BR    R11                      Redrive update logic
          ENDIF ,

```

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-----+-----1-----+-----2-----+-----3-----+-----4-----+-----5-----+-----6-----+-----7-----+-----8
* ### CODE FOR TDB EXAMINATION GOES HERE
* All diagnostics have been performed on the TDB
* Redrive the program logic with processor assist for this child
* process.
        AHI    R5,1                Increase trns abort count
        CIJNL  R5,20,DISTRAN       Failed 20 times?
        PPA    R5,0,1              Perform proc assist
        BR     R4                  Redrive transaction
*
* *****
* Child process static storage
* *****
*
STATICAREA      DC 0D'0'           align storage
S_ISGENQ_S      ISGENQ MF=(L,S_,0D)
S_ISGENQ_S_L    EQU *-S_ISGENQ_S
S_ENQ_MAJOR_NAME DC CL8'TBMSTR'
S_ENQ_MINOR_NAME DC CL(32)'TBCHILD Enqueue Res'
S_ENQ_RNAME_LENGTH DC AL1(L'S_ENQ_MINOR_NAME)
S_WAIT_TIME DC F'1'             1 timer interval - approx 26 micro secs
*
* WTO macros
*
S_WTO_SUPPORT_Y WTO 'Running in transactional mode',MF=L
S_WTO_SUPPORT_YL EQU *-S_WTO_SUPPORT_Y
S_WTO_SUPPORT_C WTO 'Running in constrained trans mode',MF=L
S_WTO_SUPPORT_CL EQU *-S_WTO_SUPPORT_C
S_WTO_SUPPORT_N WTO 'Running in non-transactional mode',MF=L
S_WTO_SUPPORT_NL EQU *-S_WTO_SUPPORT_N
S_WTO_TABORT WTO 'Transaction aborted',MF=L
S_WTO_TABORT_L EQU *-S_WTO_TABORT
S_WTO_DIS_TRAN WTO 'Disabling transactional mode',MF=L
S_WTO_DIS_TRAN_L EQU *-S_WTO_DIS_TRAN
S_WTO_TEND WTO 'Transaction ended',MF=L
S_WTO_TEND_L EQU *-S_WTO_TEND
S_WTO_TNCS WTO 'Non-constrained transaction started',MF=L
S_WTO_TNCS_L EQU *-S_WTO_TNCS
S_WTO_TNCE WTO 'Non-constrained transaction ended',MF=L
S_WTO_TNCE_L EQU *-S_WTO_TNCE
S_WTO_TCS WTO 'Con-constrained transaction started',MF=L
S_WTO_TCS_L EQU *-S_WTO_TCS
S_WTO_TCE WTO 'Con-constrained transaction ended',MF=L
S_WTO_TCE_L EQU *-S_WTO_TCE
S_WTO_ENDED WTO 'Child process ended',MF=L
S_WTO_ENDED_L EQU *-S_WTO_ENDED
S_WTO_WAITS WTO 'Child process starting wait',MF=L
S_WTO_WAITS_L EQU *-S_WTO_WAITS
S_WTO_WAITE WTO 'Child process finished wait',MF=L
S_WTO_WAITE_L EQU *-S_WTO_WAITE
*
        LTORG
STATICAREA_E    DC 0D'0'           align storage - end of static

```

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-----+-----1-----+-----2-----+-----3-----+-----4-----+-----5-----+-----6-----+-----7-----+-----8
*
* *****
* Child process working storage
* *****
*
WORKING_STORAGE DSECT ,           Program dynamic area
WS_EYE          DC CL(C_EYE_LENGTH)' '   Eye catcher
WS_SAVEA        DC 18F'0'                save area
WS_ISGENQ_S     DC 0D,XL(S_ISGENQ_S_L)'00' enq
WS_ENQ_TOKEN    DC CL(32)' '            enq token
WS_CELL_COUNT_1 DC F'0'
WS_CELL_COUNT_2 DC F'0'
WS_CELL_1P_GET  DC (C_POOL1_COUNT)A(0)   pointer to cell 1
WS_CELL_2P_GET  DC (C_POOL2_COUNT)A(0)   pointer to cell 2
WS_CELL_1P_REL  DC (C_POOL1_COUNT)A(0)   pointer to cell 1
WS_CELL_2P_REL  DC (C_POOL2_COUNT)A(0)   pointer to cell 2
WS_CELL_1HEAD   DC A(0)                  head of chain
WS_CELL_2HEAD   DC A(0)                  head of chain
WS_CELL_1TAIL   DC A(0)                  temp tail
WS_CELL_2TAIL   DC A(0)                  temp tail
WS_PSATOLD      DC A(0)                  tcb
WS_WTO_ENDED    DC 0D,XL(S_WTO_ENDED_L)'00'
WS_LOOP_CNTR    DC F'0'
*
* *****
* TRANSACTIONAL-EXECUTION FACILITY DATA AREA
* *****
*
WS_WTO_SUPPORT_Y DC 0D,XL(S_WTO_SUPPORT_YL)'00'
WS_WTO_SUPPORT_C DC 0D,XL(S_WTO_SUPPORT_CL)'00'
WS_WTO_SUPPORT_N DC 0D,XL(S_WTO_SUPPORT_NL)'00'
WS_WTO_TABORT     DC 0D,XL(S_WTO_TABORT_L)'00'
WS_WTO_DIS_TRAN   DC 0D,XL(S_WTO_DIS_TRAN_L)'00'
WS_WTO_TEND       DC 0D,XL(S_WTO_TEND_L)'00'
WS_WTO_TNCS       DC 0D,XL(S_WTO_TNCS_L)'00'
WS_WTO_TNCE       DC 0D,XL(S_WTO_TNCE_L)'00'
WS_WTO_TCS        DC 0D,XL(S_WTO_TCS_L)'00'
WS_WTO_TCE        DC 0D,XL(S_WTO_TCE_L)'00'
WS_WTO_WAITS      DC 0D,XL(S_WTO_WAITS_L)'00'
WS_WTO_WAITE      DC 0D,XL(S_WTO_WAITE_L)'00'
WS_TRAN_EYE_S     DC CL(C_EYE_LENGTH)' '   Transaction support eyecatcher
WS_TRAN_SUPPORT   DC C'N'                 Transaction support
*
WS_TRAN_TDB       DC 0D                    Transaction Diagnostic block           +0
WS_TDB_FORM       DC XL.8'0'               TDB FORMAT                             +0
WS_TDB_FLAGS      DC XL.8'0'               TDB FLAGS                               +0
                  DC 4C'0'                 ! Reserved                             +0
WS_TDB_TND        DC XL.16'0'              TDB TND (Nesting depth)                +0
WS_TDB_TAC        DC XL8'0'                TDB TRANSACTION ABORT CODE              +8
WS_TDB_CONFL      DC XL8'0'                TDB CONFLICT LOCATION                   +16
WS_TDB_ATIA       DC XL8'0'                TDB ABORTED TRANS INSTR ADDR            +24
WS_TDB_EAID       DC XL.8'0'               TDB EAID                                 +32
WS_TDB_DXC        DC XL.8'0'               TDB DXC                                  +32
                  DC 2X'0'                 ! Reserved                                 +32
WS_TDB_PI_ID      DC XL.32'0'              TDB PROGRAM INTERRUPTION ID              +32

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-----1-----2-----3-----4-----5-----6-----7-----8
WS_TDB_EXCP_ID DC XL8'0'      TDB TRANSLATION EXCEPTION ID      +40
WS_TDB_BEAL    DC XL8'0'      TDB BREAKING EVENT ADDRESS      +48
                DC 72X'0'      ! Reserved                      +56
WS_TDB_GPR     DC 120X'0'     TDB GPRs                      +128
C_TDB_LEN      EQU *-WS_TRAN_TDB
*
WS_TRAN_EYE_E  DC CL(C_EYE_LENGTH)' '    Transaction support eyecatcher
*
* *****
* END OF TRANSACTIONAL-EXECUTION FACILITY DATA AREA
* *****
*
WS_EYE_E       DC CL(C_EYE_LENGTH)' '    Eye catcher
                DC 0D          align end
WS_LENGTH      EQU *-WORKING_STORAGE WS Length
*
* Other DSECTs and COPYBOOKS
*
                COPY  TBDATA
                IHAPSA DSECT=YES,LIST=YES
                END    TBCHILD

```