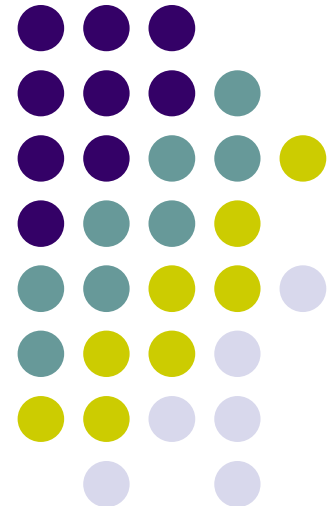
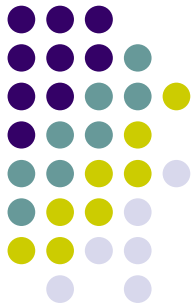


Converting your Language Environment® C/C++ Applications to XPLink for 64ZSP04587USE

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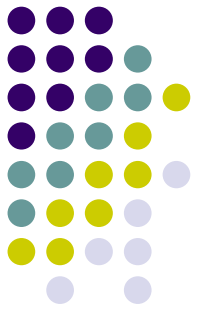
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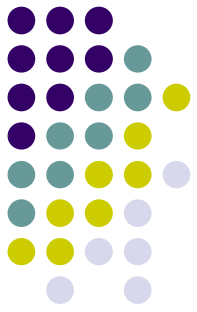
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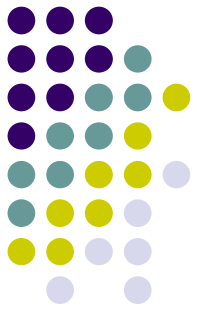
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Background

- New workloads on OS/390® and z/OS® consist largely of applications written on other platforms where function calls were "free".
- Some applications measure ~25% of execution time spent in function call overhead!
 - These are the most serious in Object-Oriented applications, where functions tend to be smaller. That is, a higher ratio of functions calls to lines of "application code".
- New, especially ported, workloads are primarily all C or C++, with little or no COBOL or Assembler.
 - This is the real target audience for XPLink.

XPLink Overview

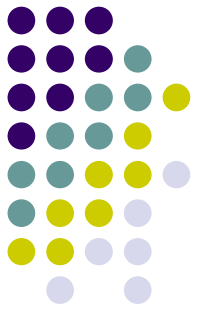


The objective of XPLink is to provide, for a specific type of application:

- Improved call linkage performance (up to 50% reduction in linkage instructions)
- Reduced function footprint in memory
- A common linkage for C/C++ (and DLLs)
- Compatibility with existing (non-XPLink) code
- No effect on existing applications

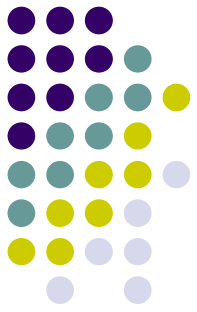
These are the characteristics expected to be exploited by 64-bit applications. Thus 64-bit is strictly XPLink!

XPLink Overview



What's Happening with Today's Linkage?

- It will probably be with us "forever".
 - At least in the existing environment
 - But perhaps not in future environments that are incompatible for other reasons (64-bit, for example, is completely XPLink)
- There is compatibility support between old and new linkages across Program Object (DLL call) boundaries.
 - At some cost in performance
 - For 31-bit only (64-bit is completely XPLink)



XPLink Overview

Did you say “improved call linkage performance”?

- YES! (Well for specific applications.) Following is a function prolog linkage comparison (the instructions that set up a new stack frame):

"Old" 31-bit f2() prolog

```
000000 47F0 F022      B      34(,r15)
000004 01C3C5C5      CEE eyecatcher
000008 00000098      DSA size
00000C 000000C0      =A(PPA1-f1)
000010 ... stack extension path
000022 90E4 D00C      STM   r14,r4,12(r13)
000026 58E0 D04C      L     r14,76(,r13)
00002A 4100 E098      LA    r0,152(,r14)
00002E 5500 C314      CL    r0,788(,r12)
000032 4130 F03A      LA    r3,58(,r15)
000036 4720 F014      BH    20(,r15)
00003A 58F0 C280      L     r15,640(,r12)
00003E 90F0 E048      STM   r15,r0,72(r14)
000042 9210 E000      MVI   0(r14),16
000046 50D0 E004      ST    r13,4(,r14)
00004A 18DE      LR    r13,r14
```

```
void f1(void) {
    f2();
    NOP flag
};

void f2(void) {
    // ...
};
```

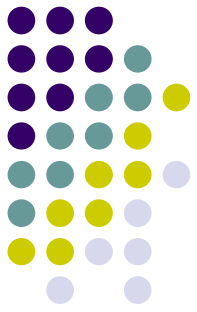
XPLink 31-bit f2() prolog

```
000000 9057 4784      STM   r5,r7,1924(r4)
000004 A74A FF80      AHI   r4,H'-128'
```

XPLink 64-bit f2() prolog

```
000000 EB57 4708 0024  STMG  r5,r7,1800(r4)
000006 A74B FF00      AGHI  r4,H'-256'
```

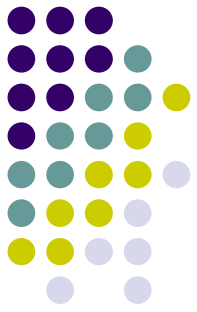
XPLink Major Differences



New calling convention

- Registers 13, 14, and 15 are just work registers
- R6 and R7 are used for linkage (ie. BASR 7,6)
 - formerly R14 and R15 were used for linkage (ie. BALR 14,15)
- R5 contains called function's own portion of WSA (its environment)
 - formerly R0 contained address of WSA, the called function had to compute the address of its own piece of Writeable Static
- No base register (R6) assumed on entry, call *may* have been made via Relative Branch
- Return register (R7) not preserved, caller cannot rely on it on return

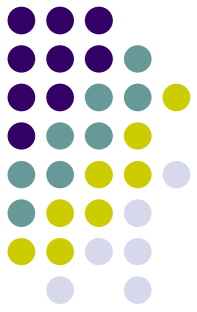
XPLink Major Differences



Improved Parameter / Return Value Passing

- Arguments passed via fixed location in caller's stack
- Arguments are directly addressable by called function
 - Also addressable by caller using same base reg as its own auto storage
- First 3 argument words passed in GPRs 1-3 (31-bit)
- First 3 argument doublewords passed in GPRs 1-3 (64-bit)
- Up to 4 floating point arguments passed in FPRs
- Remaining arguments passed in storage
- Return value is in GPR 3 (extended value in R2 and R1)

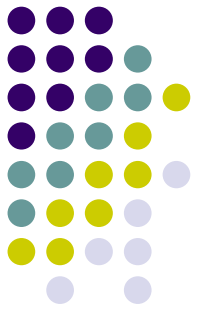
XPLink Major Differences



New Stack Layout

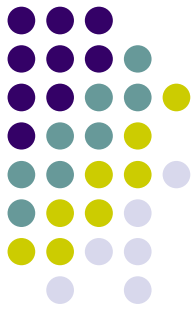
- Grows towards lower addresses
 - Called function knows how to adjust stack pointer; it doesn't have to be passed from caller (LE's NAB)
 - Beware of stack overruns with downward growing stack!
 - For example, buffer overflow beyond array's bounds extends automatic storage and overwrites calling function's stack frame.

XPLink Major Differences



New Stack Overflow Detection

- No explicit test (well.. usually) for stack overflow
 - Based on guard page size
 - Functions with large stack frames still need explicit test
 - **In 64-bit large is immense**
- Called function stores into new stack frame, storage protection used to determine if stack needs extending
 - Prolog consists of updating stack pointer and saving registers there - registers are stored in called function's stack frame, not caller's



XPLink Major Differences

Where Have all the Instructions Gone?

"Old" 31-bit f2() prolog

0000	47F0	F022	B	34(,r15)
0004	01C3C5C5			CEE eyecatcher
0008	00000098			DSA size
000C	000000C0			=A(PPA1-f1)
0010	... stack extension path			
0022	90E4	D00C	STM	r14,r4,12(r13)
0026	58E0	D04C	L	r14,76(,r13)
002A	4100	E098	LA	r0,152(,r14)
002E	5500	C314	CL	r0,788(,r12)
0032	4130	F03A	LA	r3,58(,r15)
0036	4720	F014	BH	20(,r15)
003A	58F0	C280	L	r15,640(,r12)
003E	90F0	E048	STM	r15,r0,72(r14)
0042	9210	E000	MVI	0(r14),16
0046	50D0	E004	ST	r13,4(,r14)
004A	18DE		LR	r13,r14

moving control information out of line

better register conventions reduce number of registers saved (7 vs 3 here)

detecting overflow with guard page

static instead of dynamic stack information (a compiler option can be used to force the saving of the backchain by increasing the range of the initial STM instruction)

downward-growing stack allows these 3 instructions to be replaced with a single AHI

XPLink 31-bit f2() prolog

0000	9057	4784	STM	r5,r7,1924(r4)
0004	A74A	FF80	AHI	r4,H'-128'

biased stack pointer allows this instead of:

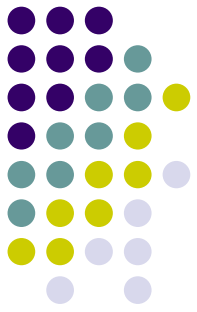
```
AHI    r4,H'-128'
*      wait for register 4 to be available
STM    r6,r7,8(r4)
```

other improvements:

- no base register
- no Library Work Area
- no stack marking

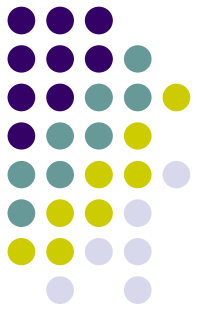
```
void f1(void) {
    f2();
    NOP flag
};

void f2(void) {
    // ...
};
```



XPLink Register Conventions

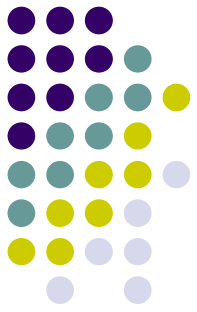
Register	XPLink Usage on Entry	XPLink Usage on Return
GPR 0	undefined	not preserved
GPR 1	1st word of argument list or undefined	not preserved or 1st word of a returned aggregate
GPR 2	2nd word of argument list or undefined	not preserved or 2nd word of returned aggregate or high half of 64-bit integer return value
GPR 3	3rd word of argument list or undefined	not preserved or 3rd word of returned aggregate or 31-bit return value
GPR 4	Address of Stack Frame minus 2048 bytes	preserved
GPR 5	Address of called function's environment or, for internal functions, containing scope's stack frame	not preserved
GPR 6	Entry point address or undefined	not preserved
GPR 7	Return address	not preserved
GPR 8-11	undefined	preserved
GPR 12	Undefined; Or for 31-bit LE-conforming applications: Pointer to (thread-specific) CAA - must be set on entry to any function	preserved (in either case)
GPR 13-15	undefined	preserved



XPLink Register Conventions

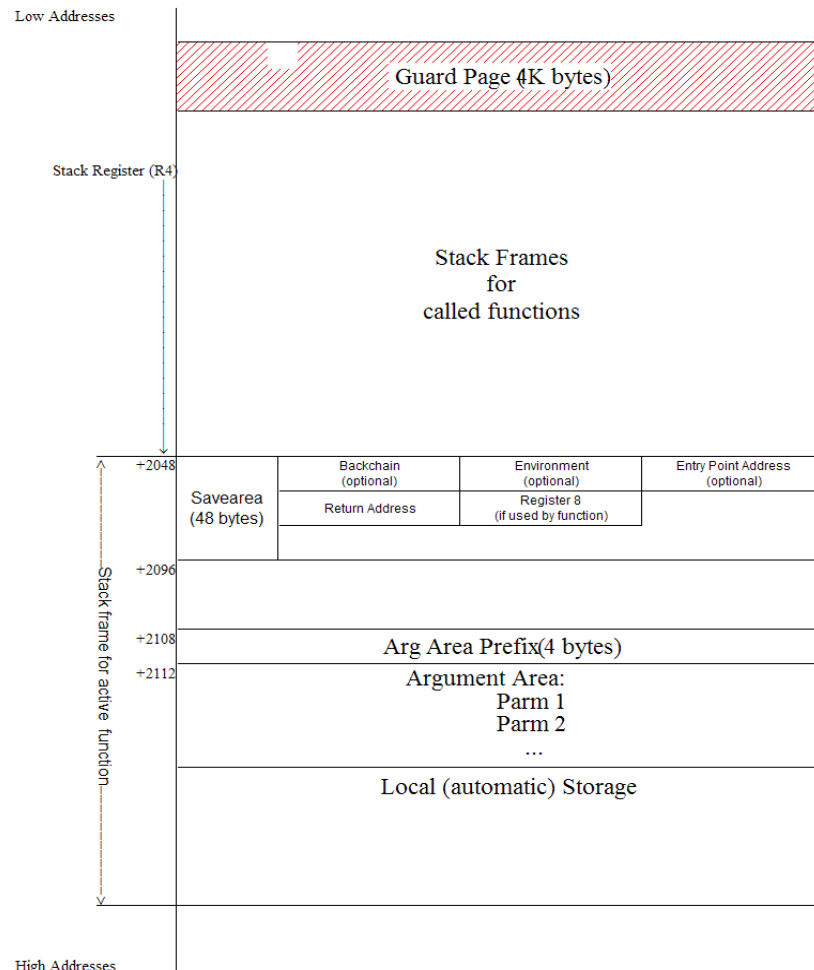
Comparing “Old” vs. XPLink Register Conventions:

Register Usage	“Old”	XPLink
Stack Ptr	R13	R4 (biased)
Return Addr	R14	R7
Entry pt on entry	R15	R6 (unless called by branch relative)
Environment	R0 (WSA)	R5
CAA Address	R12	- R12 in 31-bit - from LAA in 64-bit
Input Parm List	Address in R1	In caller's DSA, first 3 ints in R1, R2, R3, float pt values in FPR0, 2, 4, 6
Return Code/value	R15	R3 (extended value in R2 and R1)
Start of callee's stack frame	Caller's NAB value	Caller's R4 minus callee's stack frame size



XPLink Function Call Example

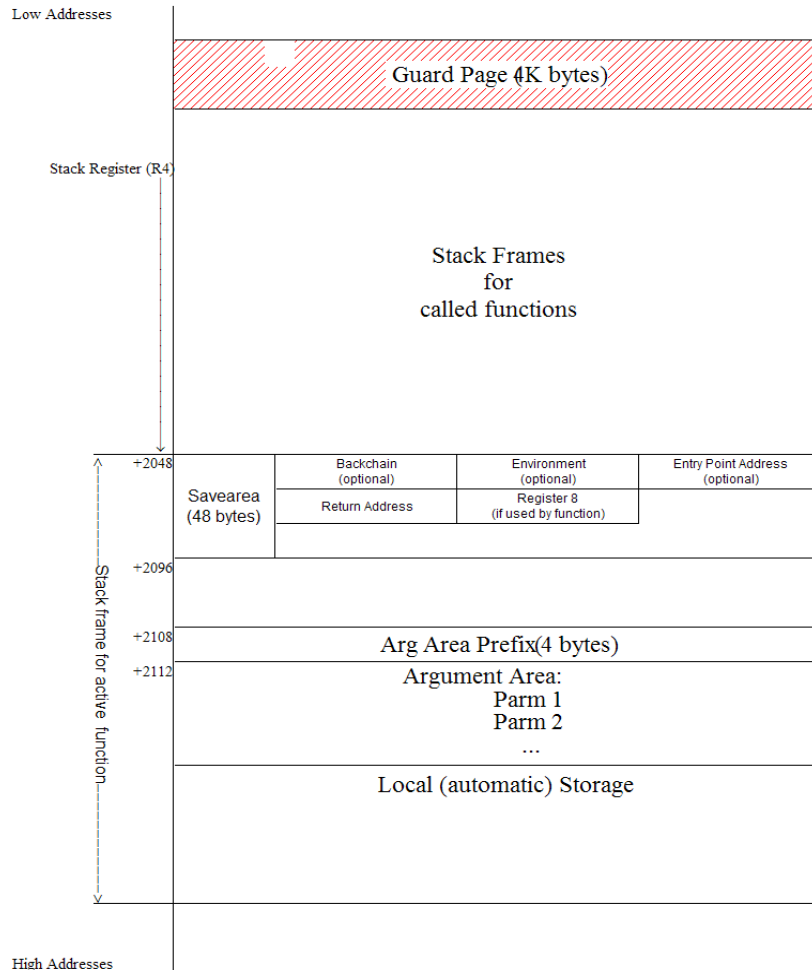
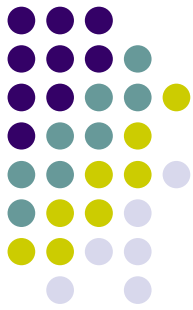
Following the XPLink stack during a function call.



```
void f1(void) {  
    f2();  
    nop flag  
};  
  
void f2(void) {  
    // ...  
    // ...  
    // ...  
};
```

1 Construct arguments in current stack frame's argument area
some arguments are passed in registers; more about this later

XPLink Function Call Example



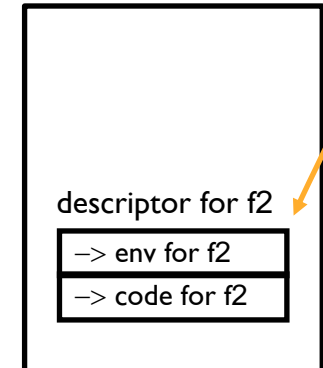
```
void f1(void) {
    f2();
    nop flag
};

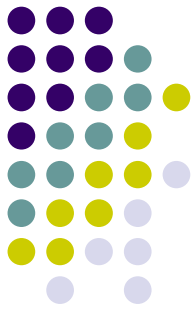
void f2(void) {
    // ...
    // ...
    // ...
};
```

```
000050 9856 5000 LM r5,r6,=A(f2)(r5,0)
000054 0D76 BASR r7,r6
000056 4700 FFFD NOP 4093(,r15)
```

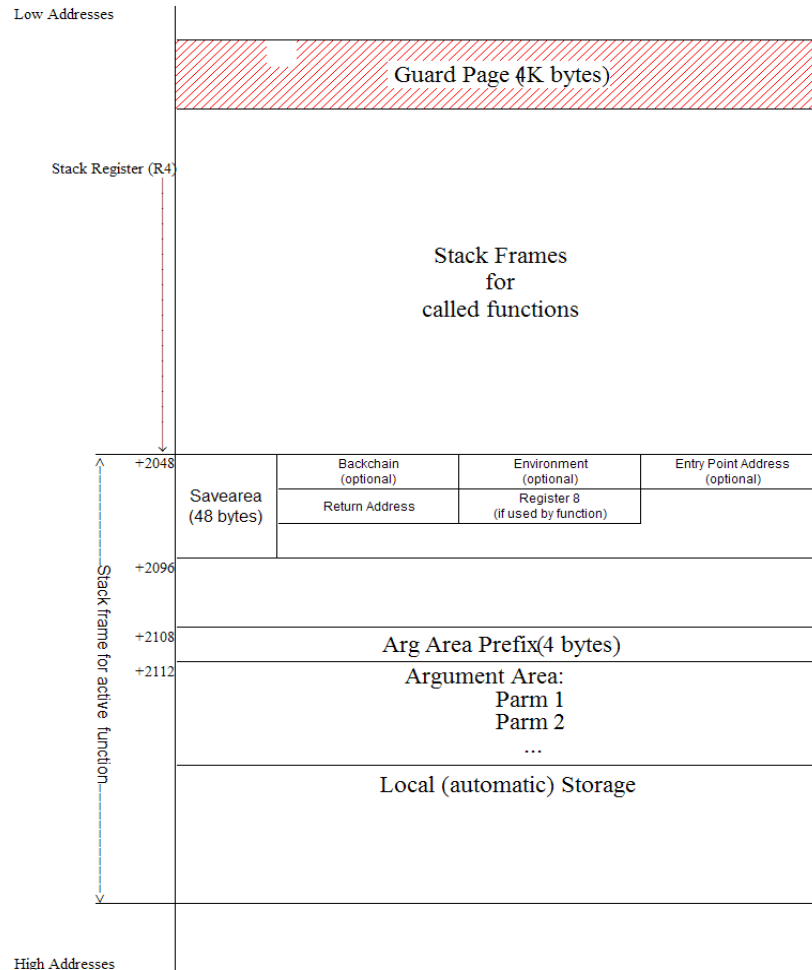
- 1 Construct arguments in current stack frame's argument area
- 2 Pick up address of function's "environment" and code from the function descriptor
the environment is an area of Writeable Static associated with the called function

Environment for f2 in WSA





XPLink Function Call Example



```

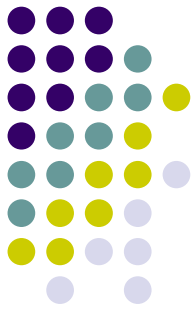
void f1(void) {
    f2();
    nop flag
};

void f2(void) {
    // ...
    // ...
    // ...
};
    
```

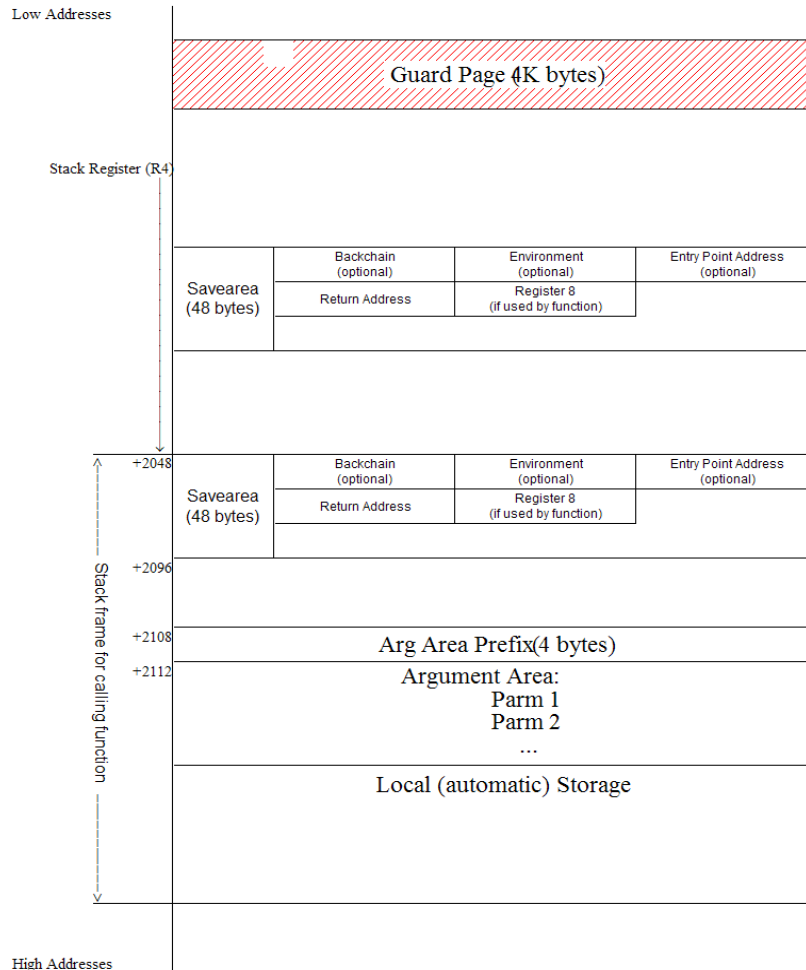
```

000050  9856  5000  LM   r5,r6,=A(f2)(r5,0)
000054  0D76                BASR r7,r6
000056  4700  FFFD                NOP  4093(,r15)
    
```

- 1 Construct arguments in current stack frame's argument area
- 2 Pick up addresses of function's "environment" and code from the function descriptor
- 3 **Link (BASR or BRAS) to called function**
 if the called function is in a DLL, the LM will pick up a handle for the called function and the address of the DLL loader, as initialized by the Binder on detecting that the called function is imported



XPLink Function Call Example



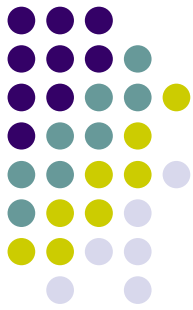
```
void f1(void) {
    f2();
    nop flag
};

void f2(void) {
    // ...
    // ...
    // ...
};
```

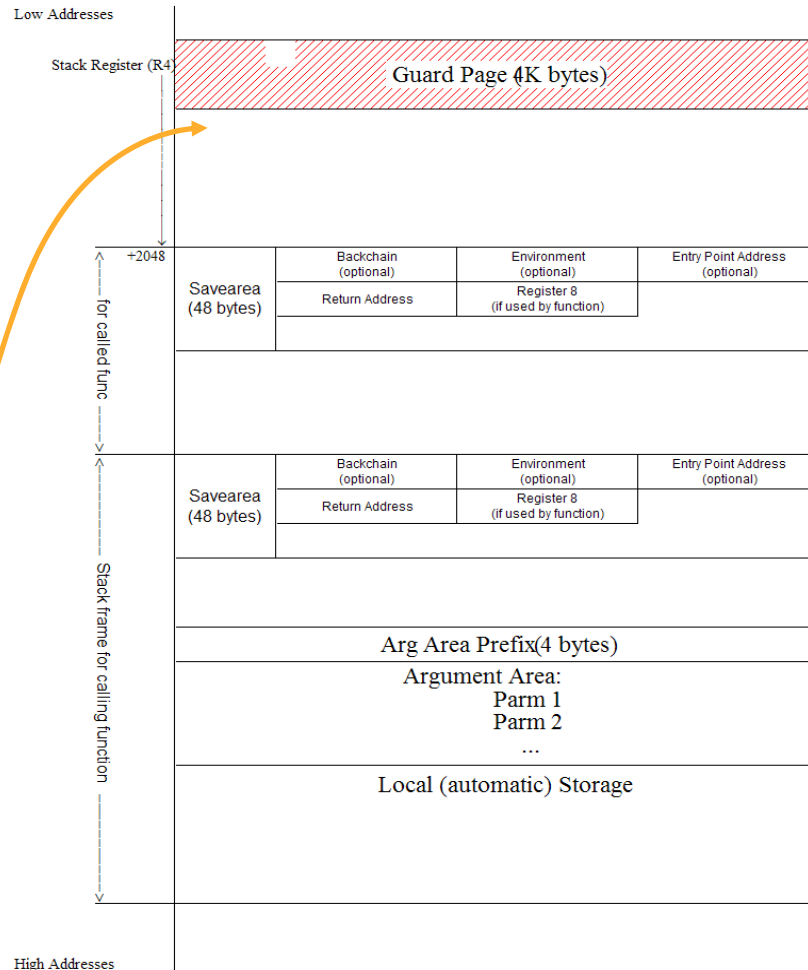
```
000048 9067 4788 STM r6,r7,1928(r4)
00004C A74A FF80 AHI r4,H'-128'
```

- 1 Construct arguments in current stack frame's argument area
- 2 Pick up addresses of function's "environment" and code from the function descriptor
- 3 Link (BASR or BRAS) to called function
- 4 **Store registers used by called function**

if the previous stack frame is close to the guard page this store (STM) may touch the guard page, causing an *interrupt that will be handled by the run time*; who, in turn, will extend the stack



XPLink Function Call Example



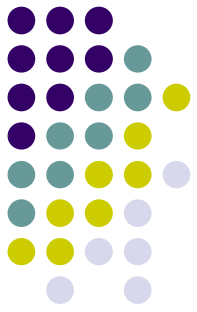
```
void f1(void) {
    f2();
    nop flag
};

void f2(void) {
    // ...
    // ...
    // ...
};
```

```
000048 9067 4788 STM r6,r7,1928(r4)
00004C A74A FF80 AHI r4,H'-128'
```

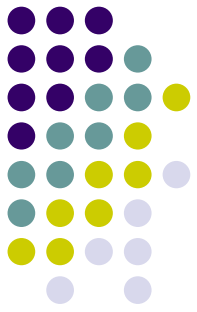
- 1 Construct arguments in current stack frame's argument area
- 2 Pick up addresses of function's "environment" and code from the function descriptor
- 3 Link (BASR or BRAS) to called function
- 4 Store registers used by called function
- 5 Update Stack Register

the stack register may (or may not) now point into the guard page; this is of no concern for this particular call



Ideal XPLink Applications

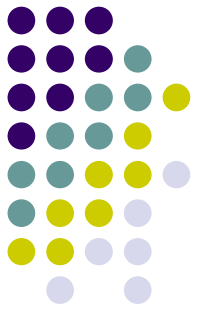
- Highly-modular with many calls to small functions
- Minimal number of calls between XPLINK and NOXPLINK-compiled functions (which requires expensive stack switching glue code) -- **not applicable to 64-bit**
 - XPLink cross-linkage support is provided in C/C++ Compiler and for High Level Assembler (via macros)
 - In general, you cannot bind XPLINK-compiled and NOXPLINK-compiled functions together in the same program object -- **not applicable to 64-bit**
- XPLink requires Binder, and the output executable must reside in a PDSE or the UNIX file system.



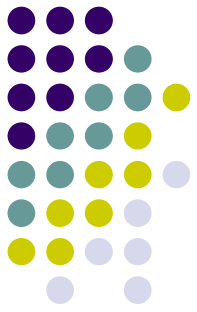
Non-Ideal XPLink Applications

- The following can degrade performance or otherwise make using XPLink unattractive:
 - **31-bit only issues**
 - Large number of cross-linkage calls between XPLink and non-XPLink functions (requires stack switching glue code)
 - In an XPLink environment the C RTL is compiled XPLINK so non-XPLink callers of C functions go through stack switching glue
 - The C RTL uses stack switching glue code internally
 - Hex Math Library requires stack switch to run on Upstack (IEEE Floating Pt Library is ok, it's XPLink)
 - **31-bit and 64-bit issue**
 - Using unsupported environment or function (more on this later)

Cross-linkage Function Calls



- Since XPLINK and NOXPLINK-compiled parts cannot be mixed in the same program object, the DLL calling mechanism is the primary method for calling between XPLink and non-XPLink
 - Also supported are `fetch()/fetchep()` and LE's CEEFETCH macro
- The following do **not** support calls to XPLINK-compiled functions:
 - COBOL Dynamic Call
 - PL/I Fetch
 - CEELOAD (i.e.. the traditional LOAD/BALR)

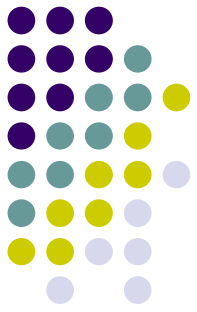


Cross-linkage Function Calls

Ok, so there is *some* support for calling non-XPLink functions statically from XPLink:

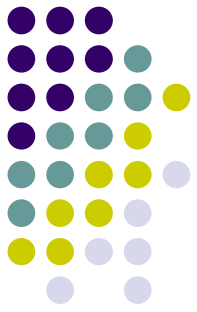
- **#pragma linkage(..., OS_NOSTACK)**
 - In **31-bit**, this is identical to OS31_NOSTACK
 - Generates direct call using OS Linkage conventions (no glue, so fast) but only 72-byte (31-bit)/144-byte (64-bit) savearea (e.g. C headers)
- **#pragma linkage(..., OS_UPSTACK) -- 31-bit only**
 - Generates call to RunOnUpStack glue so called function gets control with OS Linkage conventions and LE-conforming stack
- The intent is to be able to call Assembler "leaf" routines to perform functions not easily done from C/C++.

Stack Switching Glue Code



- Calls between XPLink and non-XPLink functions require LE to insert "glue code" that will:
 - switch between the upward and downward growing stacks
 - adjust parameter list formats
 - non-XPLink uses R1 pointing to list of parameters (or parm addresses)
 - XPLink passes parameters in general and floating point registers

Stack Switching Glue Code

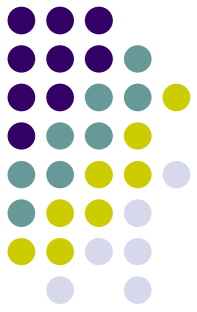


- Services provided to allow non-XPLink and XPLink routines to run on the correct upward- or downward-growing stack:

- ▶ **CEEVROND (RunOnDownstack)** -- calls XPLINK-compiled function from non-XPLink caller
- ▶ **CEEVRONU (RunOnUpstack)** -- calls NOXPLINK-compiled function from XPLink caller
- ▶ **CEEVH2OS (XPLink-to-OSLinkage)** -- calls non-XPLink function from XPLink callers using OS Linkage conventions

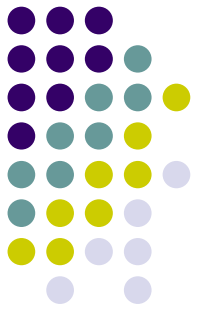
Refer to the z/OS Language Environment Vendor Interfaces book for details on these CWIs (Compiler-Writer Interfaces).

Unsupported Environments



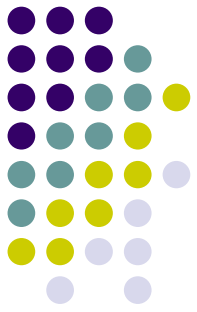
- ~~CICS®~~ (added in TS 3.1)
- ~~DB2®~~ stored procedures (EXEC SQL is allowed)
- ~~IMS™~~ transactions (calls to ctdli() allowed)
- ~~PIP1~~ (added in z/OS v1r3)
- PIC1
- LRR
- 64-bit additional limitations
 - AMODE-31 and non-LE conforming applications
 - Nested enclaves
- AMODE-24 and non-LE conforming applications
- Child nested enclave must match parent enclave's XPLINK run-time option
- CEEBXITA and CEEBINT User Exits cannot be coded as XPLINK functions

Building an XPLink Application



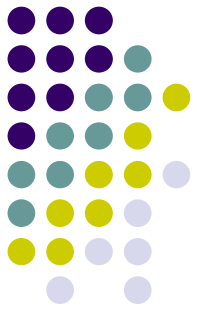
	Non-XPLINK...	XPLINK...
SYSLIB static libraries	SCEELKED SCEELKEX SCEE OBJ SCEECPP	SCEEBND2 SCEEBIND
Dynamic Link Library (DLL) side decks	None for LE	In SCEELIB: 31-bit: <ul style="list-style-type: none">▶ CELHS003 (C RTL)▶ CELHS001 (LE AWIs)▶ CELHSCPP (C++) 64-bit: <ul style="list-style-type: none">▶ CELQS003 (C/LE RTL)▶ CELQSCPP (C++)

Building an XPLink Application



- SCEEBND2 is a new LE data set containing XPLINK-compiled static routines ("stubs")
 - There are only a few
 - This data set can only be used with XPLINK applications (SCEELKED, etc. are non-XPLINK only)
- SCEELIB is a new LE data set containing LE DLL side decks
 - For XPLINK applications, the C RTL *is* a DLL

Building an XPLink Application

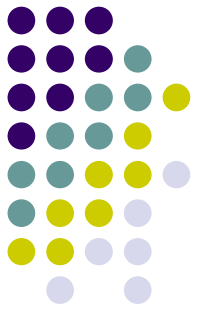


XPLINK Compile Option

NOXPLINK | XPLINK (optional suboptions)

- XPLINK(BACKCHAIN | NOBACKCHAIN)
 - With BACKCHAIN, STM instruction in prolog begins with register 4 to provide an explicit link between stack frames. This is not necessary for tools like CEEDUMP and slows down the prolog code.
- XPLINK(STOREARGS | NOSTOREARGS)
 - With STOREARGS, compiler inserts extra code after prolog to explicitly store parameter registers into argument area in caller's DSA.

Building an XPLink Application

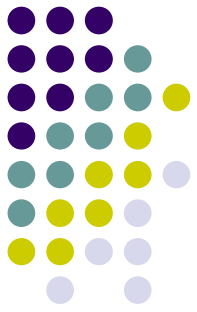


XPLINK Compile Option

NOXPLINK | XPLINK (optional suboptions)

- XPLINK(OSCALL(Downstack | Upstack | Nostack))
 - Alters default behavior of #pragma linkage(..., OS)
- XPLINK(NOGUARD | GUARD)
 - NOGUARD will generate an explicit check of the stack floor in the prolog code
- XPLINK(NOCALLBACK | CALLBACK) **z/OS R2**
 - CALLBACK will allow non-XPLink function pointers or descriptors to be correctly used in an XPLink program. `__callback` qualifier is preferred.

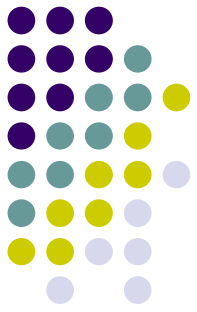
Building an XPLink Application



c89 changes

- The XPLINK compile option can be specified as:
 - `-Wc,xplink` (`-Wc,goff` is forced)
 - `-Wc,lp64` (`-Wc,xplink,arch(5+)` forced, `-Wc,float(ieee)` defaulted)
 - Object files are still Fixed 80, but now in GOFF format
- A new XPLINK linkedit option is also required (this option is *not* passed to the binder):
 - `-WI,xplink`
 - `-WI,lp64`
 - Tells c89 to use SCEEBND2 and SCEELIB data sets
 - Forces binder options `DYNAM=DLL` and `CASE=MIXED` (required for calls to C RTL and other DLLs)

Building an XPLink Application



c89 simple example

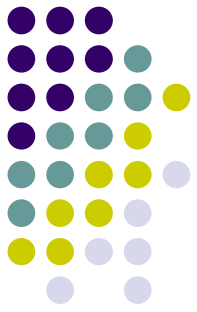
- XPLink "Hello World" example:

```
c89 -o HelloWorld -Wc,xplink -Wl,xplink HelloWorld.c
```

- 64-bit "Hello World" example:

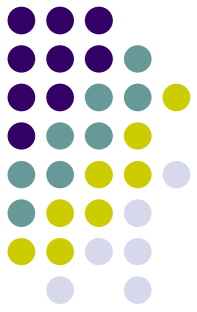
```
c89 -o HelloWorld -Wc,lp64 -Wl,lp64 HelloWorld.c
```


Running an XPLink Application



- **XPLink Requires that both LE run-time libraries are available at execution time:**
 - SCEERUN
 - SCEERUN2
 - It's a PDSE (required by XPLink)
 - Contains XPLink versions of C RTL, locales and converters, etc.
 - 👉 31-bit: CELHV003, CELHDCPP (C++)
 - 👉 64-bit: CELQLIB (combined C/LE RTL), CELQDCPP (C++)

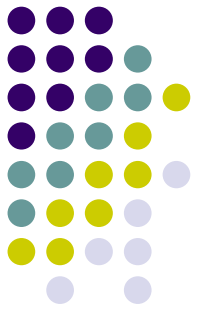
Running an XPLink Application



New LE Run-Time Option - 31-bit

- XPLINK(ON|OFF)
 - XPLINK(OFF) is the default
 - If main() is compiled XPLINK, then the XPLINK run-time option is forced ON
 - If main() is compiled NOXPLINK *but* calls an XPLINK-compiled function, then XPLINK(ON) must be specified, otherwise error message CEE3555S will be generated and the application is terminated
 - Cannot be specified in CEEDOPT as a system installation default, the XPLINK run-time option must be specified on an application by application basis (when needed)

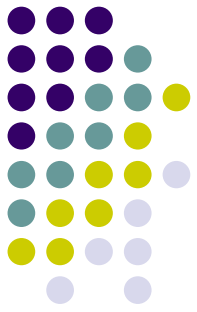
Running an XPLink Application



Changed LE Run-Time Options – 31-bit

- **STACK** (usinit_size, usinc_size, ANY|BELOW, KEEP|FREE, dsinit_size, dsinc_size)
 - **STACK suboptions:**
 - upstack initial size, upstack increment size
 - upstack location (ANY | BELOW)
 - **location ANY forced when XPLINK(ON) in effect**
 - duration (KEEP | FREE)
 - downstack initial size <-- new**
 - downstack increment size <-- new**
 - Downstack sizes do not include storage for guard page
 - Downstack not allocated in XPLINK(OFF) environment

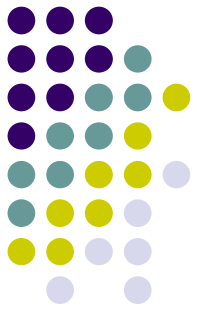
Running an XPLink Application



Changed LE Run-Time Options – 31-bit

- **THREADSTACK** (usinit_size, usinc_size, ANY|BELOW, KEEP|FREE, dsinit_size, dsinc_size)
 - **THREADSTACK** suboptions:
 - upstack initial size, upstack increment size
 - upstack location (ANY | BELOW)
 - **location ANY forced when XPLINK(ON) in effect**
 - duration (KEEP | FREE)
 - downstack initial size** <-- new
 - downstack increment size** <-- new

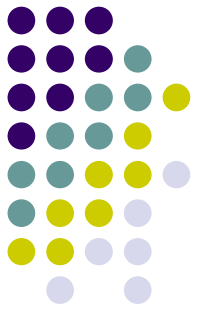
Running an XPLink Application



Changed LE Run-Time Options – 31-bit

- THREADSTACK continued...
 - Downstack sizes do not include storage for guard page
 - Downstack not allocated in XPLINK(OFF) environment
 - ☞ THREADSTACK option replaces the NONIPTSTACK and NONONIPTSTACK options (which are still accepted for compatibility)

Running an XPLink Application



Changed LE Run-Time Options – 31-bit

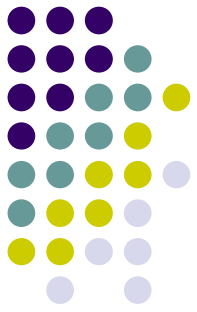
- ALL31

- When the XPLINK(ON) run-time option is in effect, the ALL31 run-time option will be forced to ON.
- No AMODE 24 routines allowed in an XPLINK(ON) environment

- RPTSTG

- Will report storage statistics for downward stack

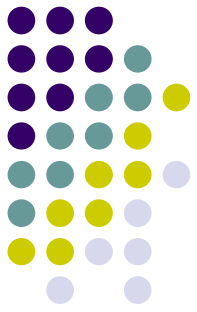
Running an XPLink Application



New LE Run-Time Options - 64-bit

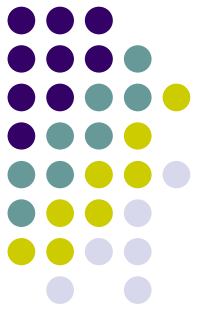
- Many unneeded options eliminated (XPLINK, ALL31)
- Many old options have **64-bit** counterparts: ...
 - **HEAPOOLS64**(OFF|ON, up to 12 *cellsize, cellcounts*)
[OFF, 8, 4000, 32, 2000, 128, 700, 256, 350, 1024, 100,
2048, 50, 3072, 50, 4096, 50, 8192, 25, 16384, 10, 32768, 5, 65536, 5]
☞ In 31-bit, HEAPPOOLS takes cell percentages rather than cell counts.
 - **HEAP64**(init, incr, KEEP|FREE,
init31, incr31, KEEP|FREE,
init24, incr24, KEEP|FREE)
[1M, 1M, KEEP, 32K, 32K, KEEP, 4K, 4K, FREE]
☞ `__malloc31()` and `__malloc24()` available for storage Below-The-Bar and Below-The-Line.

Running an XPLink Application



New LE Run-Time Options - 64-bit

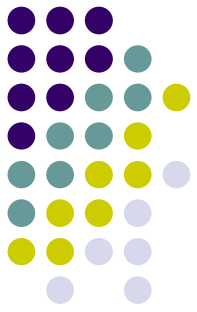
- Many old options have **64-bit** counterparts:
 - **STACK64**(init,incr,max)
[1M,1M,128M]
 - **THREADSTACK64**(OFF|ON,init,incr,max)
[OFF,1M,1M,128M]
- ☞ Stack is *always contiguous* (not so in 31-bit).
- ☞ Current stack size identified by guard page.
- ☞ On stack overflow (exception) guard page is moved down.



Compiler-writer Interfaces

The following CWIs are new for XPLink:

- Documented in LE Vendor Interfaces
- Declared in `<edcwccwi.h>` (in SCEESAMP data set)
 - `__dsa_prev()`
 - Takes as input the address and format of "current" DSA
 - Returns address of previous (logical or physical) DSA and its format
 - Call it in a loop to unwind the stack



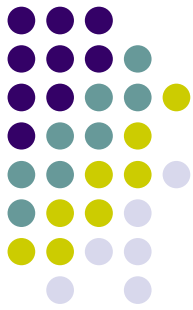
Compiler-writer Interfaces

- `__ep_find()`
 - Takes as input a DSA address and format
 - Returns the address of the entry point of the function owning the input DSA

- `__bldxfd()` - **31-bit only**
 - XPLink environment only
 - Takes a function pointer (entry point) of unknown linkage as input
 - Returns a "function pointer" that can be called by all linkage types

Appendix A

Sample Generated XPLink Code



There are also new XPLink-style entry points and Program Prolog Areas (PPAs)

```

00001 |      *  #include <stdio.h>
00002 |      *
00003 |      *  main() {

000020 |      @1L0    DS    0D          XPLink entrypoint marker
000020 | 00C300C5    =F'12779717'      '.C.E.E.1' eyecatcher
000024 | 00C500F1    =F'12910833'      (x'F1' == entry pt marker)
000028 | FFFFFFFE0   =F'-32'           Offset to XPLINK-style PPA1
00002C | 00000080    =F'128'          DSA size
000000 |      00003 |      main    DS    0D          Function entry point
000000 | 9057  4784  00003 |      STM    r5,r7,1924(r4)
000004 | A74A  FF80  00003 |      AHI    r4,H'-128'
000008 |      End of Prolog

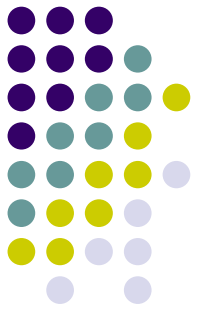
000008 | 5810  4804  00004 |      *      printf("Hello world\n");
00000C | 9856  1010  00004 |      L      r1,#Save_ADA_Ptr_1(,r4,2052)
000010 | 0D76      00004 |      LM     r5,r6,=A(printf)(r1,16)
000012 | 4700  0003  00004 |      BASR  r7,r6
000016 | 4130  0000  00004 |      NOP   3
00001A |      00005 |      *      }
00001A |      00005 |      LA    r3,0
00001A |      00005 |      @1L1   DS    0H

00001A |      Start of Epilog
00001A | 5870  480C  00005 |      L      r7,2060(,r4)
00001E | 4140  4080  00005 |      LA    r4,128(,r4)
000022 | 07F7      00005 |      BR    r7

```

Appendix B

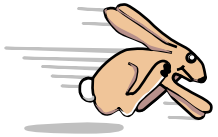
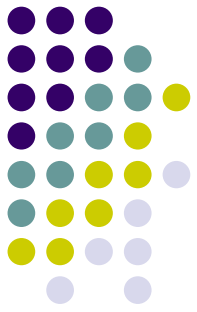
Main XPLink and 64-bit Publications of Interest



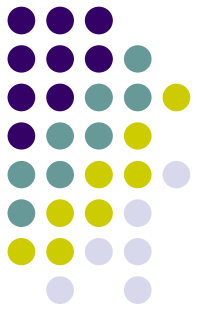
- LE Programming Guide (SA22-7561) has a chapter on developing XPLink applications.
- LE Programming Guide for 64-bit (SA22-7569)
- LE Vendor Interfaces (SA22-7568) has new detailed description, new CWIs, and "all" 3 LE-conforming linkages:
 - Standard LE linkage (includes COBOL, PL/I, etc)
 - C++ Fastlink
 - XPLINK
- LE Writing Interlanguage Communication Applications (SA22-7563)
- LE Debugging Guide (GA22-7560)
- XPLINK Redbooks™ (<http://publib-b.boulder.ibm.com/abstracts/sg245991.html>)
- 64-bit Redpaper (<http://publib-b.boulder.ibm.com/abstracts/redp9110.html>)
- C/C++ for z/OS books updated too

Appendix C

Performance Measurements - 31-bit



Reference Materials

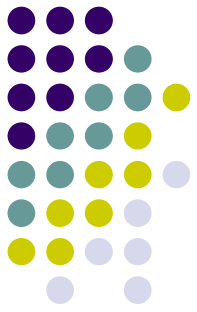


XPLink Performance Redbook

- Measurements made over summer 2000
- SG24-5991

www.redbooks.ibm.com/redbooks/SG245991.html

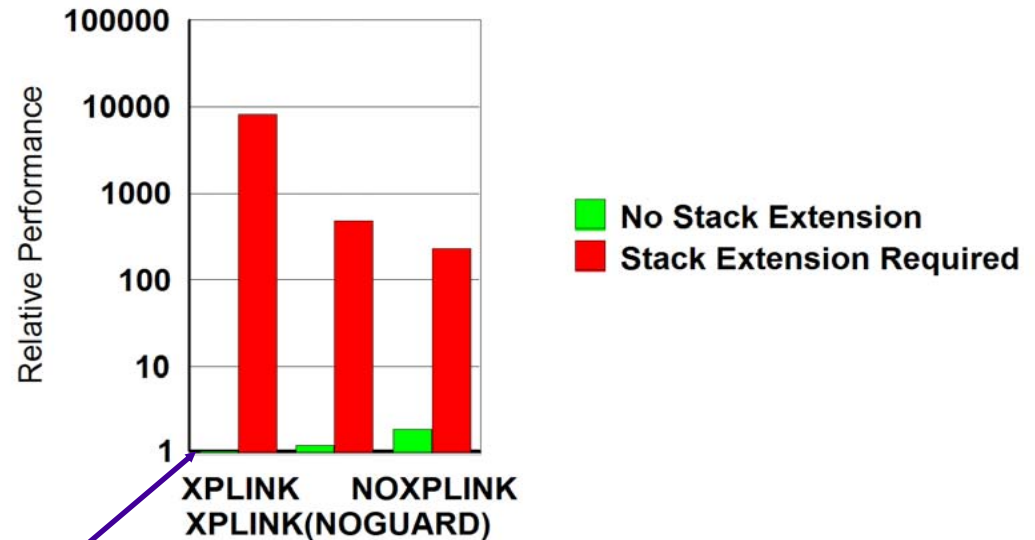
- Measurements were made on shared systems in Toronto (a development system) and Poughkeepsie (the ITSO system)
- Results were generally repeatable within 1-2%
- Highlights are reported here, details are in the Redbook



The Importance of Storage Tuning

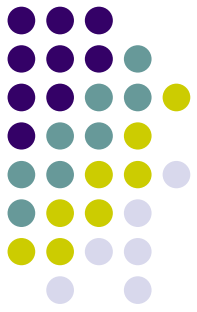
- Stack overflow detection is by program check
 - Improperly-tuned stack allocation can cause disastrous performance
 - use XPLINK(NOGUARD) in portions of the application where stack growth is unpredictable

Cost of Improper Storage Tuning



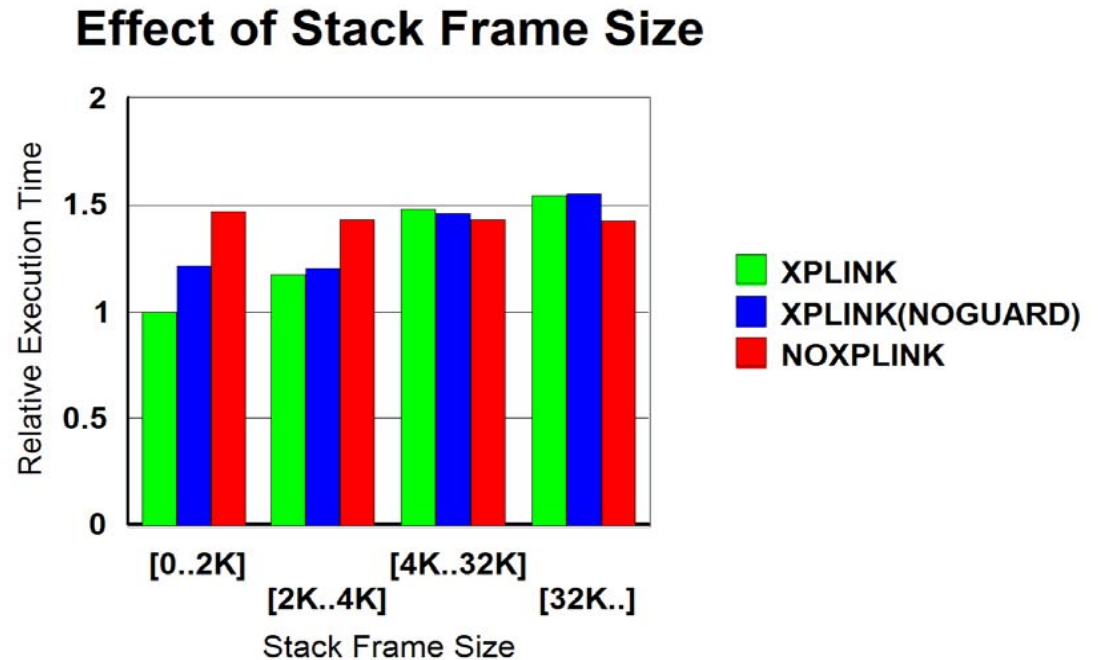
Baseline (1) is the XPLINK test running with a "proper" initial stack allocation

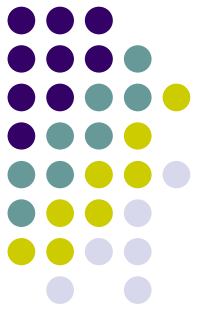
Note the logarithmic scale



Effect of Stack Frame Size

- XPLink is optimized for small stack frames (that is, small amounts of automatic storage)

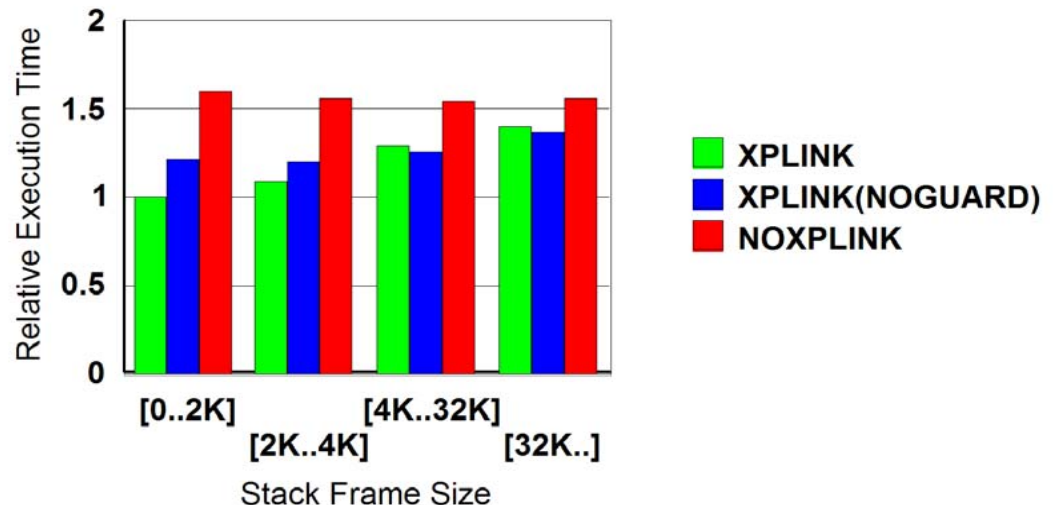


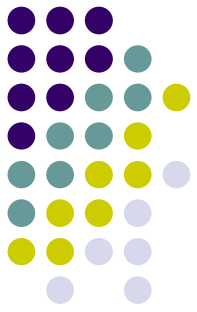


Effect of Number of Parameters

- XPLinkfunction prologs change with the number of function arguments
- Fewer arguments gives better code-generation opportunities
- The worst case (>32K local storage, 1 parameter) is better than non-XPLink

Effect of Stack Frame Size
Single Parameter

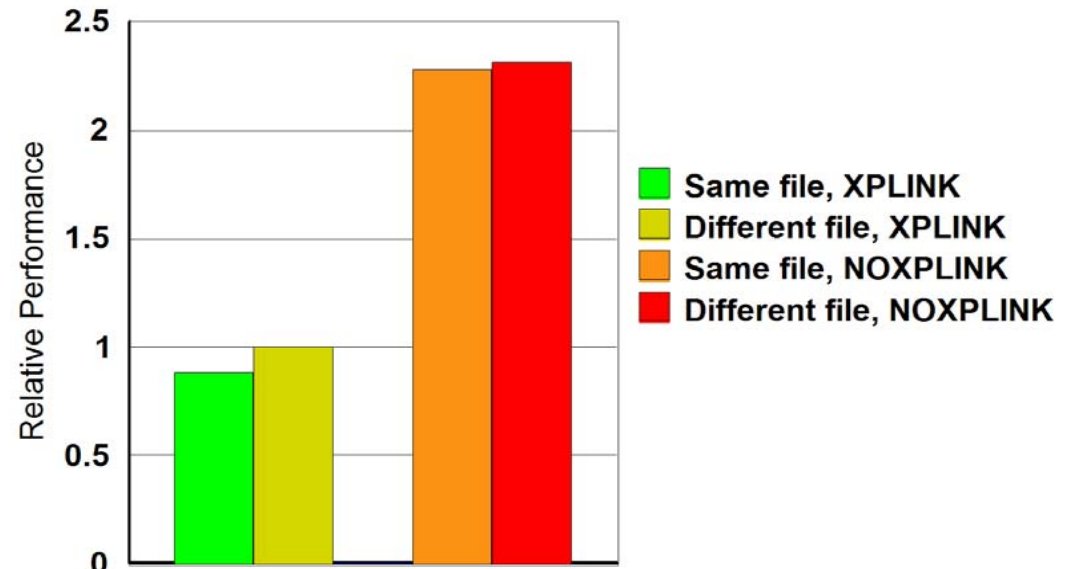




Calling Within a Compilation Unit

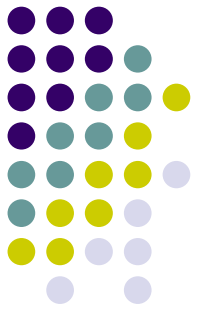
- Calls are often faster when made to a function in the same compilation unit
- Environment pointer (WSA pointer for NOXPLINK) is often the same
- Can be called with relative branch instructions in XPLink
 - ▶ function is entered with no base register

Effect of Calling Within Compilation Unit



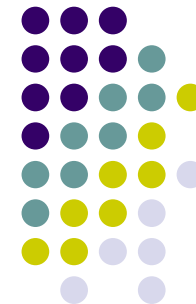
The XPLINK advantage from being within the same compilation (12%) is more pronounced than with NOXPLink (2%)

Mixing XPLink with a COBOL Application

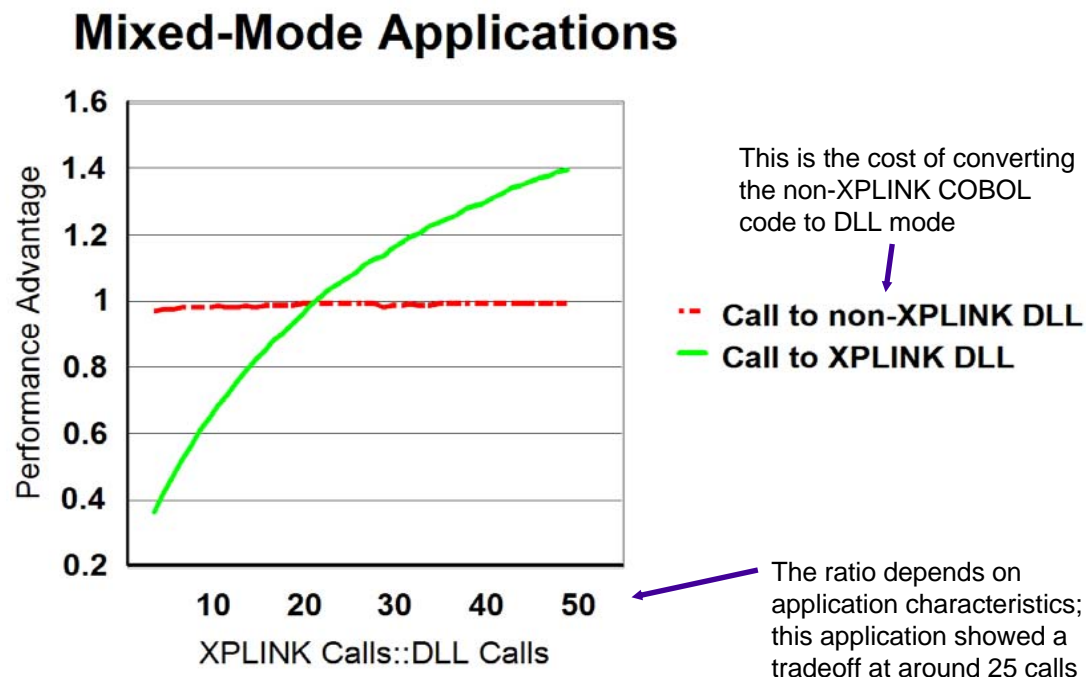


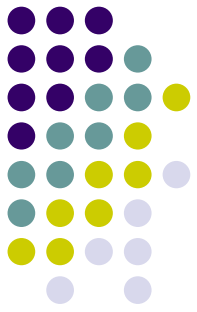
- COBOL does not support XPLink
- Separate the XPLink (C/C++) code from the non-XPLink code
- Put XPLink code into a DLL

Mixing XPLink with a COBOL Application



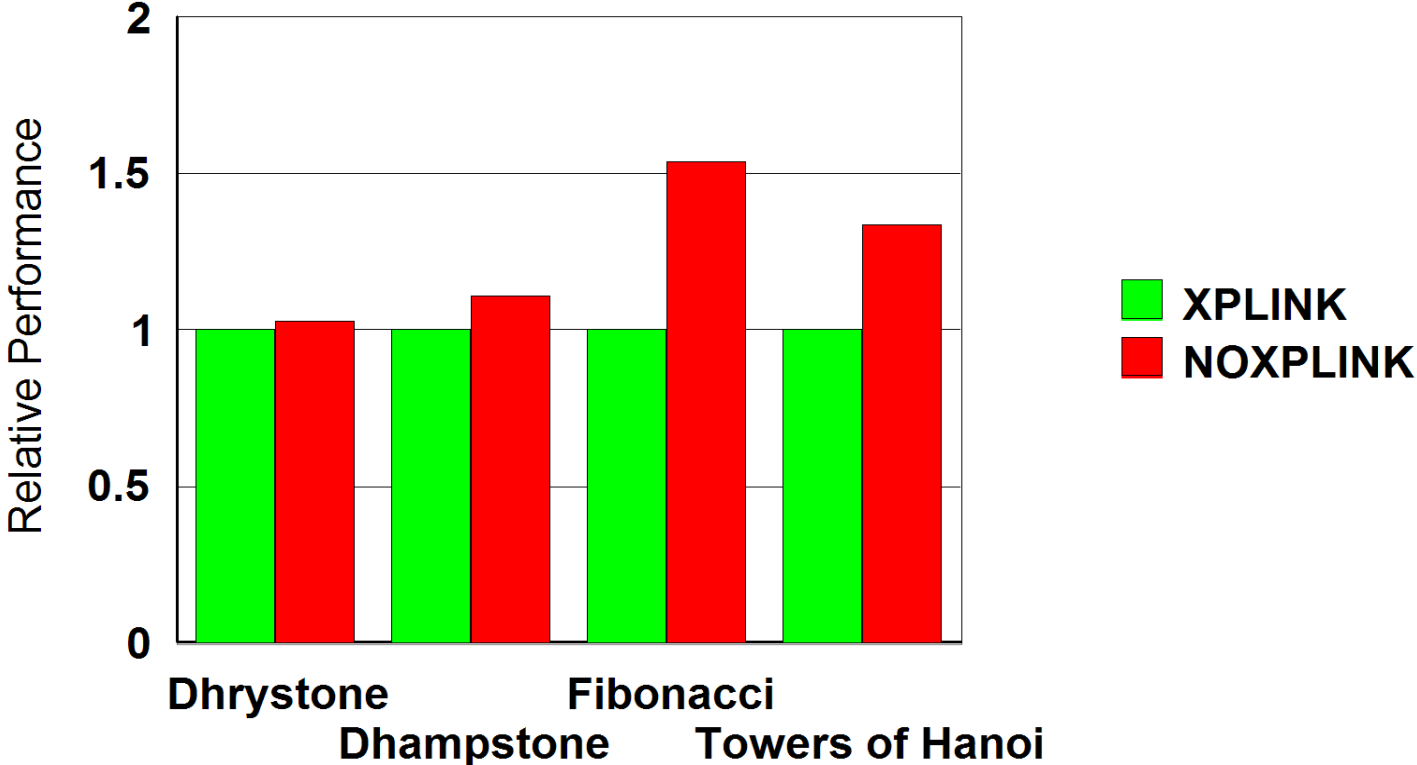
- Overall application performance depends on the number of calls *inside* the DLL for every call *into* the DLL
- This is typical of the performance characteristics expected from a C/C++ DLL written for use with a COBOL DLL.

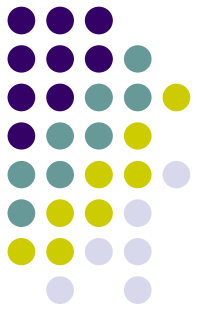




Industry Benchmarks

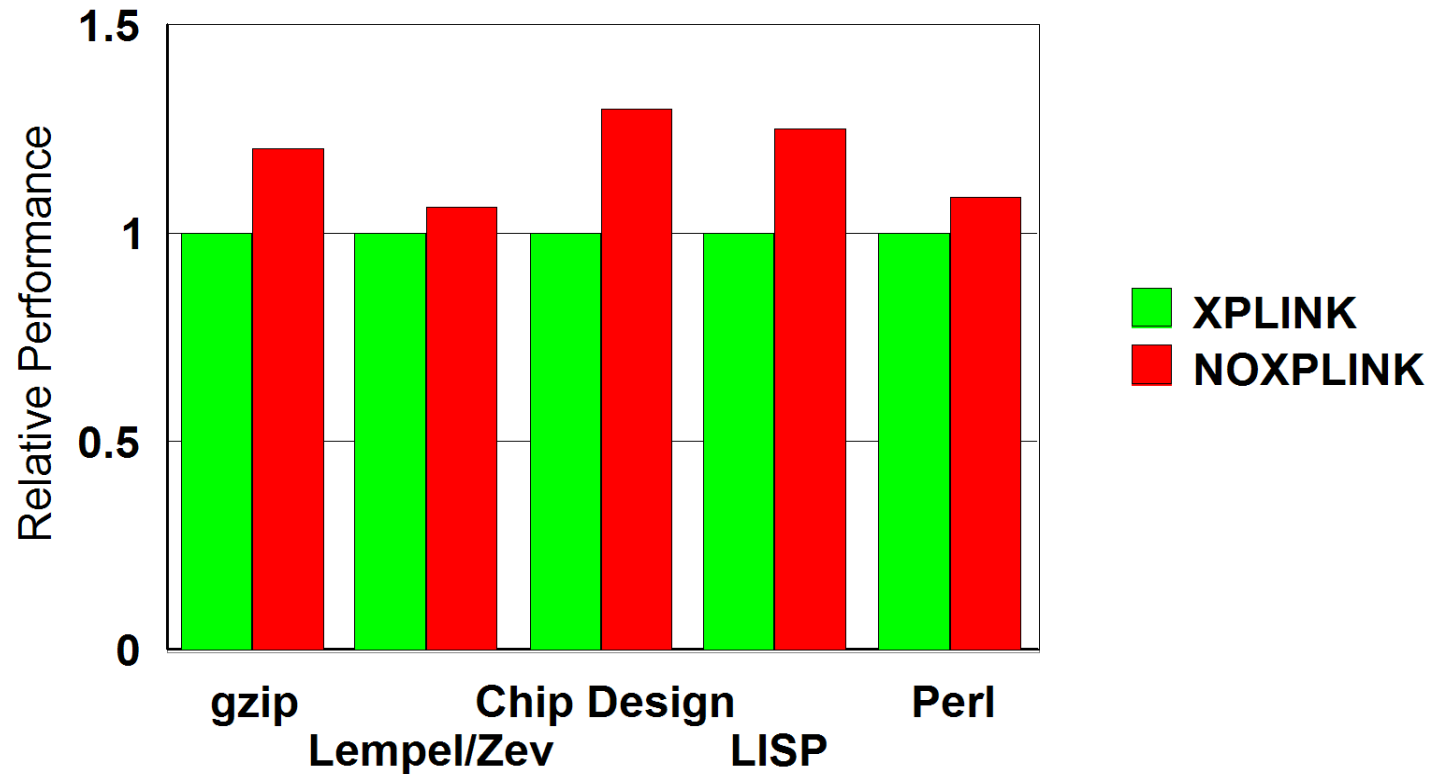
Industry Benchmarks



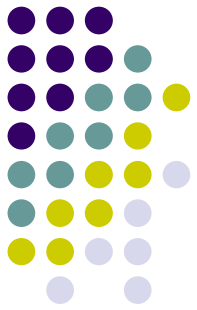


CPU Intensive Benchmarks

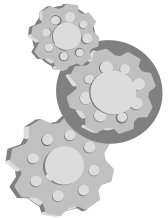
CPU Intensive Benchmarks



Appendix D

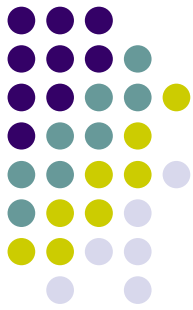


XPLink More Details



Reference Materials

Prolog Comparison for Large Automatic Storage



- Functions with large automatic storage clearly do not get the same performance advantage with XPLink

```

000060 47F0 F022      B    34(,r15)
000082 90E4 D00C      STM  r14,r4,12(r13)
000086 58E0 D04C      L    r14,76(,r13)
00008A 5800 F008      L    r0,8(,r15)
00008E 1E0E          ALR  r0,r14
000090 5500 C314      CL  r0,788(,r12)
000094 4130 F03C      LA  r3,60(,r15)
000098 4720 F014      BH  stack_extender
00009C 58F0 C280      L    r15,640(,r12)
0000A0 90F0 E048      STM  r15,r0,72(r14)
0000A4 9210 E000      MVI 0(r14),16
0000A8 50D0 E004      ST  r13,4(,r14)
0000AC 18DE          LR  r13,r14
    
```

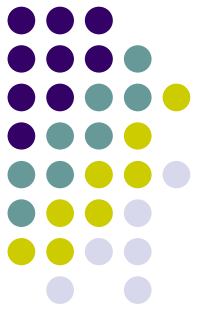
```

void f1(void) {
    f2(1,2,3);
    NOP flag
};

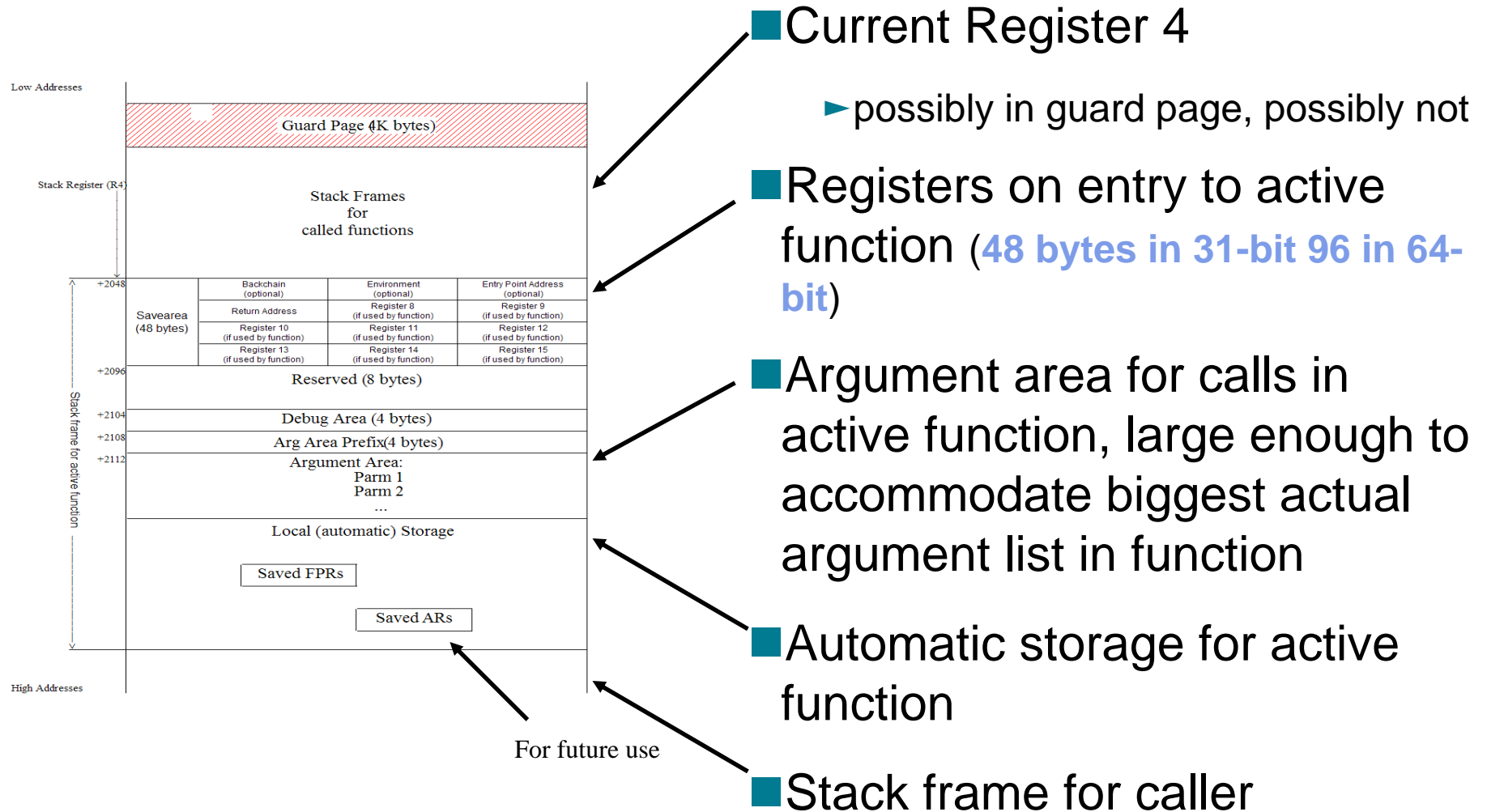
void f2(int i,int j, int k)
{
    // huge local storage
    // requirements
    // ...
    // ...
};
    
```

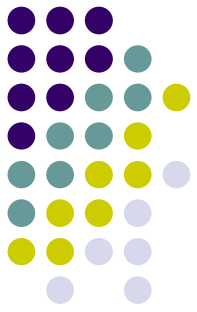
```

000048 9023 4844      STM  r2,r3,2116(r4)
00004C 1804          LR  r0,r4
00004E 0D20          BASR r2,0
000050 A72A 0034      AHI  r2,H'52'
000054 5A40 2000      A    r4,0(,r2)
000058 5940 C364      C    r4,868(,r12)
00005C A744 0022      JL  stack_extender
000060 9058 4804      STM  r5,r8,2052(r4)
000064 5000 4800      ST  r0,2048(,r4)
000068 1882          LR  r8,r2
00006A 1820          LR  r2,r0
00006C 5820 2844      L    r2,2116(,r2)
    
```

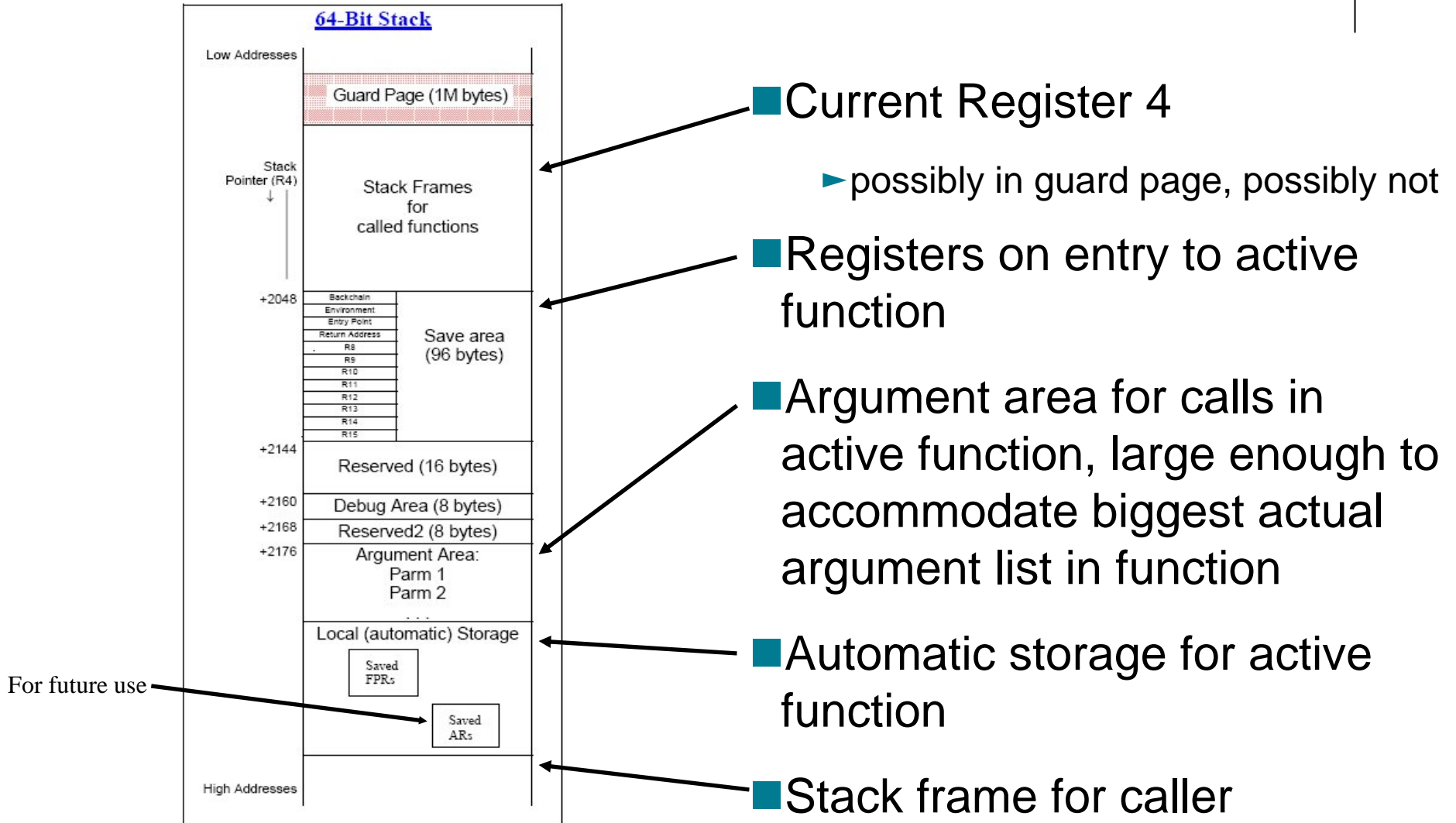



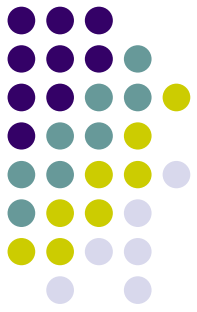
Stack Layout Detail - 31-bit



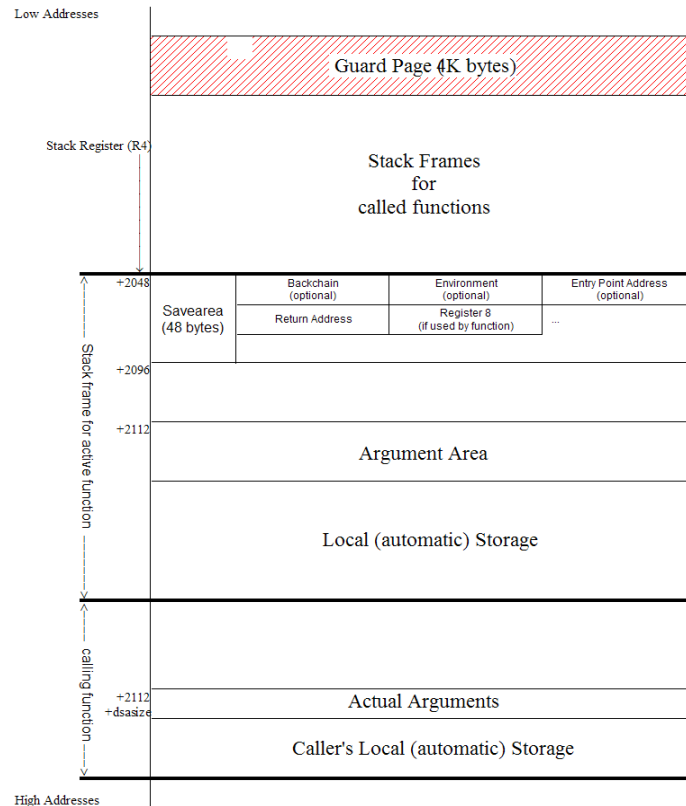


Stack Layout Detail - 64-bit

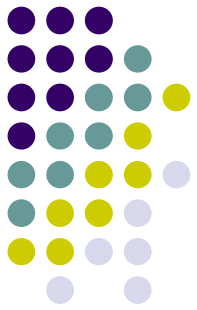




Argument Addressability



- Active Stack Frame
- Actual Arguments to Active function, built by caller in caller's argument area. Addressable at:
 - 31-bit** ($R4+2112+active\ stack\ frame\ size$)
 - 64-bit** ($R4+2176+active\ stack\ frame\ size$)



Entry Point Marker

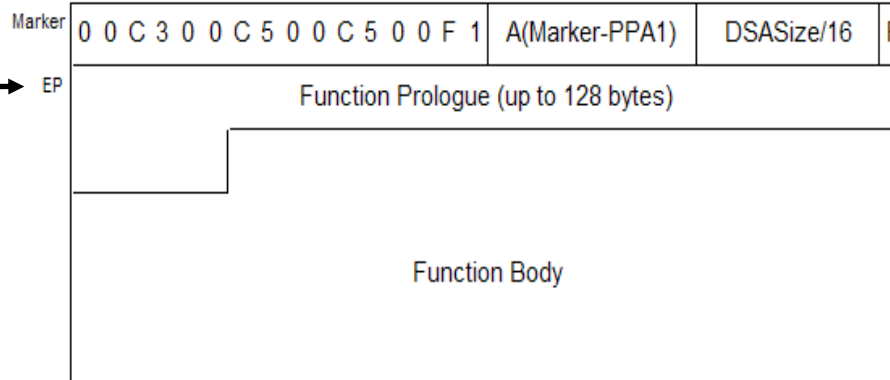
■ Entry Point

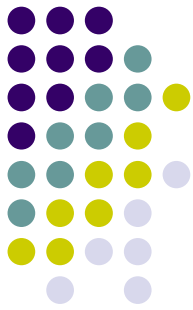
- doubleword aligned

■ Entry Point Marker

- 16 bytes before entry point, doubleword aligned
- Shows up in dump as **.C.E.E.1**

■ PPA1 Locator





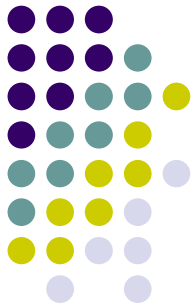
New PPA1 format

+0	Version 0x02	CEL Signature 0xCE	Saved GPR Mask	
+4	Signed offset to PPA2 from start of PPA1			
+8	PPA1 Flags 1	PPA1 Flags 2	PPA1 Flags 3	PPA1 Flags 4
+12 0x0C	Length/4 of Params		Length/2 of Prologue	Alloca Reg
+16 0x10	Offset/2 to Stack Pointer Update			
	Length of Code			

+0	State Variable Locator		PPA1 Flags 3 Bit 0
+0	Argument Area Length		PPA1 Flags 3 Bit 1
+0	FPR mask	AR mask	PPA1 Flags 3 Bit 2 or 3
+0	Floating Point Register Save Area Locator		PPA1 Flags 3 Bit 2
+0	Access Register Save Area Locator		PPA1 Flags 3 Bit 3
+0	PPA1 Member Word		PPA1 Flags 3 Bit 4
+0	Block Debug Info (PPA3) address		PPA1 Flags 3 Bit 5
+0	<u>Interface Mapping Flags</u>		PPA1 Flags 3 Bit 6
+0	Link- age	Return Value Adjust	
+0	Parameter Mapping Flags		
+4	Java Method Locator Table (MLT)		PPA1 Flags 3 Bit 7
+0	Length of Name	Name of Function	PPA1 Flags 4 Bit 7
+4			
+8			
...			

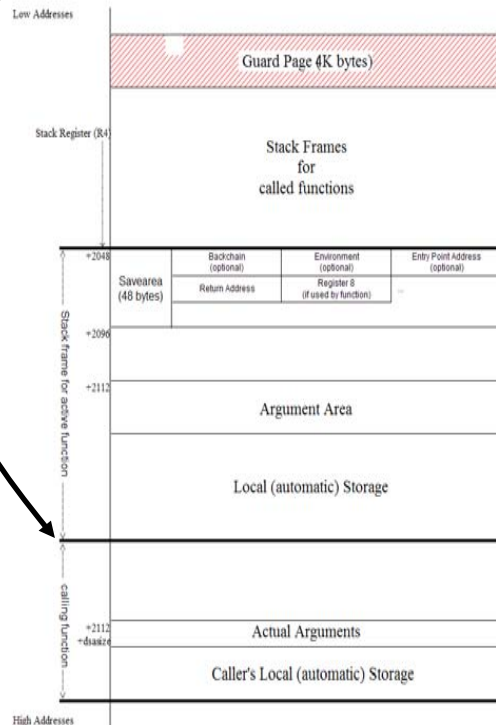
"Locators" are 4-bit register numbers followed by 28-bit offsets. Add the contents of the specified register and the specified offset to get to the target of the locator

- Fixed portion
- Optional Fields, presence indicated by flags in fixed portion



Stack Walking

The Stack:

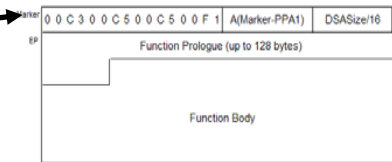


```

void f1(void) {
    f2();
    nop flag
};

void f2(void) {
    // ...
    //
    // ...
};
    
```

The Code:

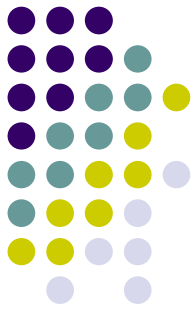


Stopped here, in dump, debugger, service routine &c

The steps described here assume we are **not** in f2()'s prolog. We can determine this by:

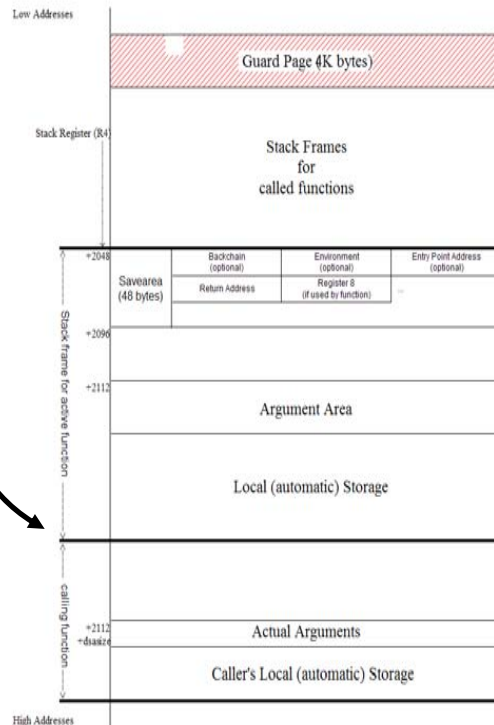
- 1 scanning backwards up to 128 bytes looking for the doubleword-aligned Entry Point Marker.
- 2 from the Entry Point Marker locating the PPA1 (as described in the following foils) and, in the PPA1, the length of the prologue and the offset of the instruction updating the stack pointer

Version 0x02	CEL Signature 0xCE	Saved GPR Mask	
Signed offset to PPA2 from start of PPA1			
PPA1 Flags 1	PPA1 Flags 2	PPA1 Flags 3	PPA1 Flags 4
Length 4 of Params		Length 2 of Prologue	Offset 2 to Stack Pointer Update
Length of Code			
State Variable Locator			PPA1 Flag 5 Bit 0
Argument Area Length			PPA1 Flag 5 Bit 1
FPR mask		AR mask	
Floating Point Register Save Area Locator			PPA1 Flag 5 Bit 2 or 3
Access Register Save Area Locator			PPA1 Flag 5 Bit 2
PPA1 Member Word			PPA1 Flag 5 Bit 4
Block Debug Info (PPA3) address			PPA1 Flag 5 Bit 5
Interface Mapping Flags			PPA1 Flag 5 Bit 6
Linkage	Return Value Adjust	Parameter Mapping Flags	
Java Method Locator Table (MLT)			PPA1 Flag 5 Bit 7
Length of Name			PPA1 Flag 4 Bit 7
Name of Function			



Stack Walking Steps

The Stack:



```
void f1(void) {
    f2();
    nop flag
};

void f2(void) {
    // ...
};
```

- 1 Pick up f2()'s return address from current stack frame
- 2 Look at call type
- 3 Pick up f2()'s entry point from current stack frame (computed in case of a relative call)
 - extract offset from the relative branch instruction just prior to the return point, add it to return address
- 4 Pick up f2()'s PPA1 offset

5 Locate f2()'s PPA1, examine f2()'s GPR Save Mask in the PPA1
 this tells us which registers were saved in f2()'s stack frame by f2()'s prologue similar rules apply to floating point registers: there's a "Floating Point Register Save Area Locator" in PPA1 which tells us where to find the FPR save area in the current stack frame, and the FPR mask which tells us which floating point registers were saved.

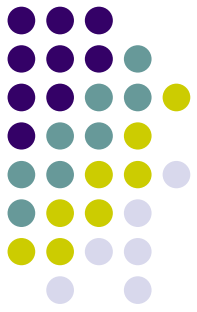
6 Store unsaved registers into f2()'s Save Area
 the first time through (while processing the stack frame for the active function) use the values actually in the registers at the time of interrupt; subsequently, use the register values stored in the previous stack frame

7 Pick up f2()'s dsasize and flags

8 Add f2()'s dsa size to current stack frame address to get f1()'s stack frame

if f2() uses alloca(), pick up the alloca register from the PPA1 and add the dsasize to that register

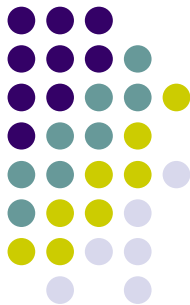
9 Repeat as required



Stack Walking Steps (continued)

- Stack structure is fully supported by
 - Debug Tool (**31-bit**)
 - dbx
 - IPCS LEDATA, described in z/OS MVS™ Interactive Problem Control System (IPCS) Commands
- Additional documentation can be found in
 - z/OS Language Environment Debugging Guide and Run-Time Messages
 - z/OS Language Environment Vendor Interfaces

Call Type Info and Call Descriptors

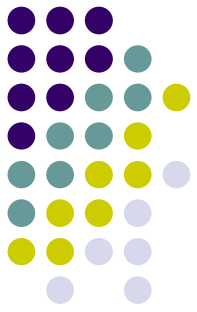


- Call Type Info: NOP(R) at call site
 - 31-bit uses NOP
 - **64-bit uses NOPR -- there is no call descriptor!**
- Used to describe return type and parameters passed in registers:

```

* f2();
  L      5,f2..env pointer
  BRAS  7,f2
  NOP   0(call type)           470t 0000
  ORG   *-2
  DC    H'(offset of EP marker or descriptor)/8'
  ...
* call descriptor (shareable)
  DS    0D
  DC    A(Signed offset to EP Marker)
  DC    AL1(return type)
  DC    AL3(parameter descriptor)
    
```

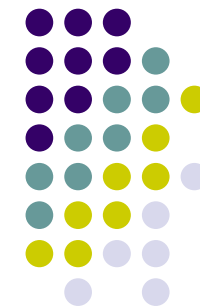
0000	Function is called with a BASR 7,6 instruction; register 6 will contain its entry point address
0001	Function is called via a BRAS 7,EP instruction; the called function does not have a base register on entry
0010	Reserved
0011	Reserved; the called function does not have a base register on entry
0100	Reserved
0101	Reserved
0110	Non-XPLink call inside XPLink function body
0111	Special linkage
1...	Reserved



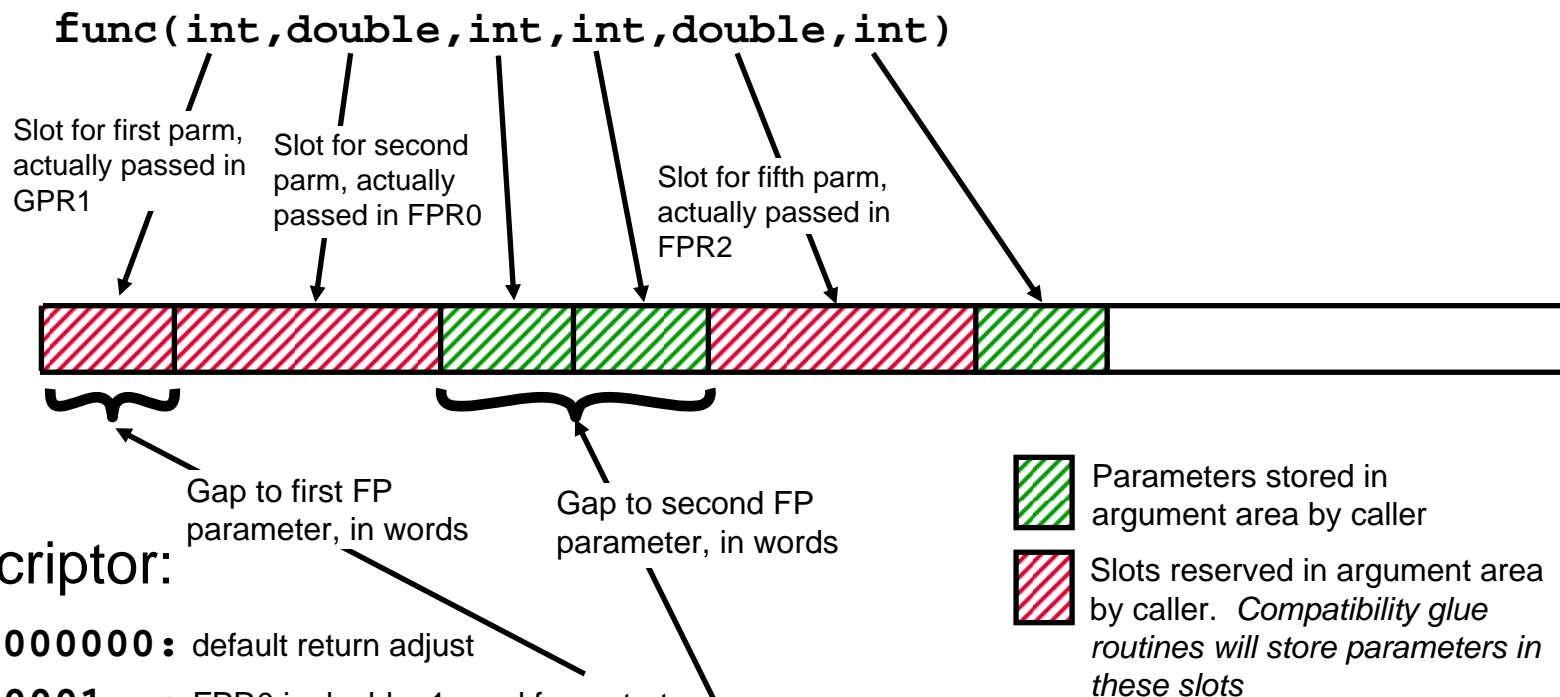
Call Descriptor - 31-bit

- Describes parameters locations for floating point parameters passed in registers
 - indicates where the "holes" in the stack-based parameter list occur
- Return adjust for return values in registers
- Solely for compatibility with non-XPLINK code

Floating Point Parameters, Example

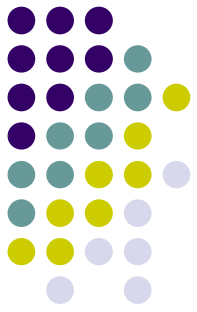


■ Doubles passed in FPR0 and FPR2



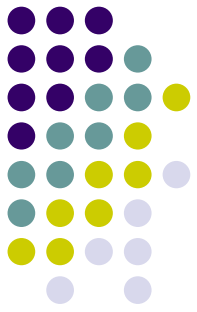
Descriptor:

- 00000000** : default return adjust
- 100001** : FPR0 is double, 1 word from start
- 100010** : FPR2 is double, 2 words from previous float
- 000000** : FPR4 not used
- 000000** : FPR6 not used



Examining Actual Arguments

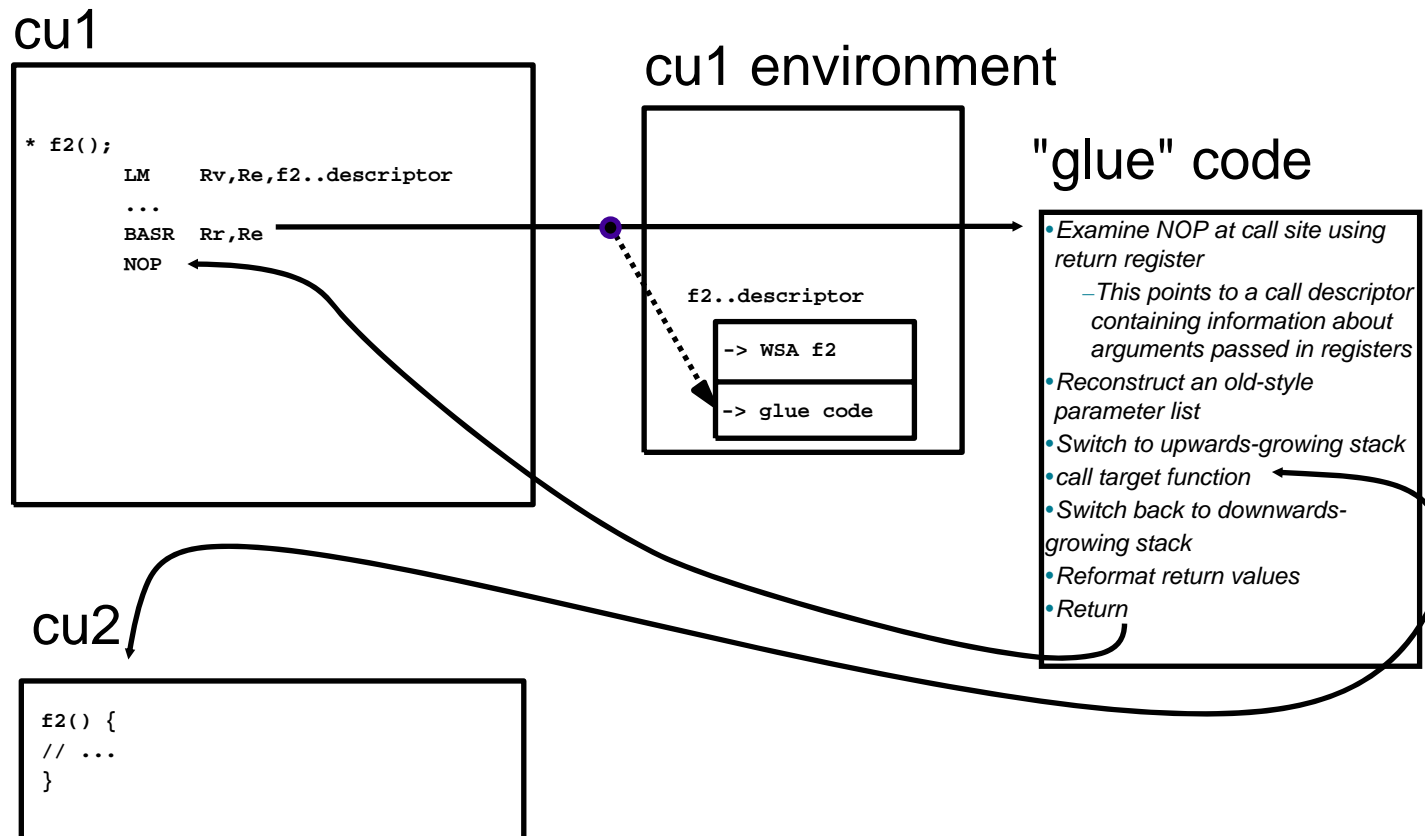
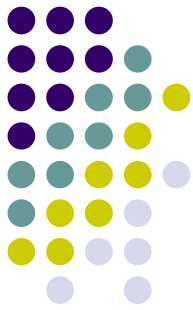
- Some arguments are passed in registers (GPRs 1-3, FPRs 0, 2, 4, 6)
- Their values may have been lost if they are no longer required
- Compiler XPLINK(STOREARGS) option forces generated code to save incoming parameters in their natural places in the caller's argument area
- Even then, they may be lost if the called function modifies its arguments



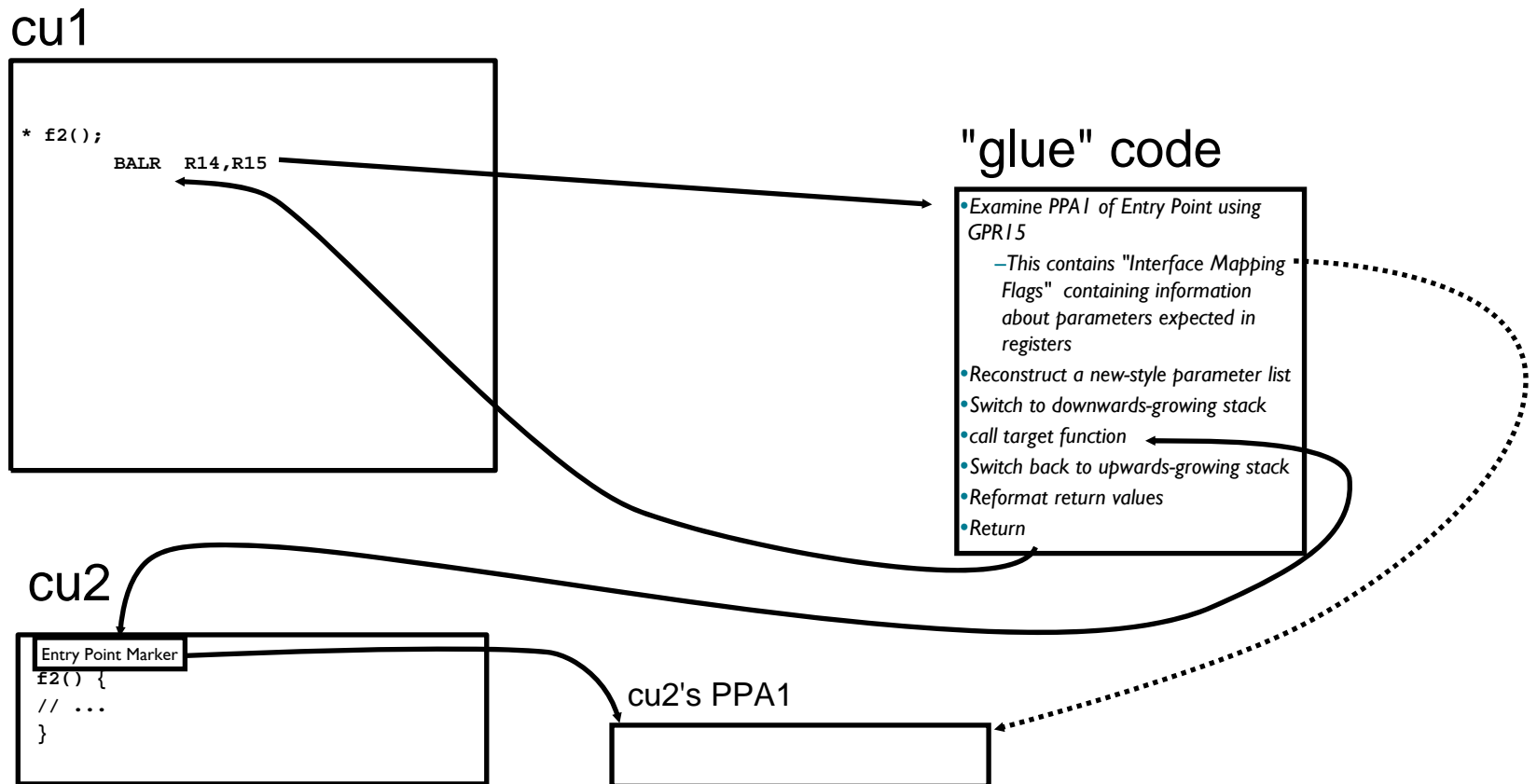
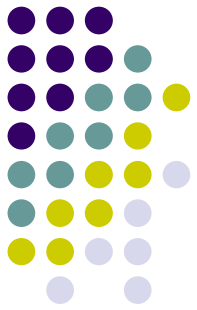
New Linkage Specifier

- New #pragma introduced to provide some low-level compatibility support
 - OS_NOSTACK, to call assembler code
 - ▶ R13 points to 18-word savearea
 - ▶ R14, R15 linkage
 - ▶ R1 points to parameters
 - ▶ No NAB, or other "Classic LE" stack artifacts
 - This is the default for linkage OS; it can be changed by a compiler option

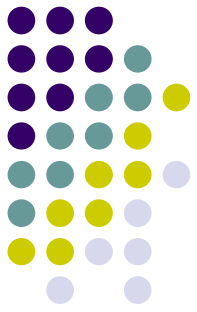
Compatibility: New Calling Old - 31-bit



Compatibility: Old Calling New - 31-bit



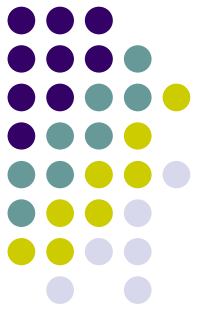
Appendix E



Debugging an XPLink Application

Reference Materials



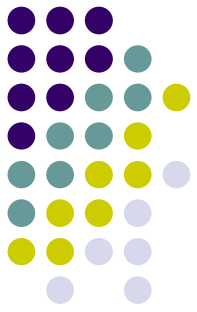


CEEDUMP Support for XPLink

- Traceback support for Up and Down stacks - **31-bit**

Traceback:

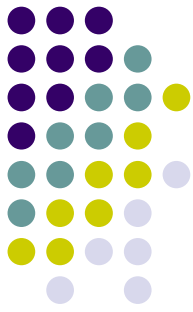
DSA Addr	Program Unit	PU Addr	PU Offset	Entry
23F91F50	CEEHDSPR	23BA0090	+000041B0	CEEHDSPR
23F914E8		23AB25E8	+0000005C	dllfunc
23F91338	CEEVRONU	23CA3348	+00000706	CEEVRONU
240316A0		23AB13D0	+00000016	main
24031720		23CA1D10	+000009A4	CEEVROND
23F910E0	EDCZHINV	23F64118	+0000009A	EDCZHINV
23F91018	CEEBBEXT	00053380	+000001A6	CEEBBEXT



CEEDUMP Support for XPLink

31-bit

- After the Traceback, the DSAs on the stack are formatted.
- Individual DSAs labeled as:
 - "UPSTACK DSA" (Non-XPLink)
 - "DOWNSTACK DSA" (XPLink)
 - "TRANSITIONAL DSA" (LE "glue")
 - CEEVRONU -- RunOnUpstack
 - CEEVROND -- RunOnDownstack



CEEDUMP Support for XPLink

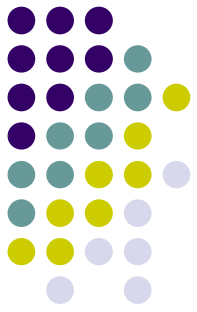
- Traceback support for 64-bit stacks

Traceback:

DSA	Entry	E Offset	Load Mod	Program Unit	Service	Status
00000001	CEEHDSP	+00000000	CELQLIB	CEEHDSP	HLE7709	Call
00000002	CEEOSIGJ	+00000956	CELQLIB	CEEOSIGJ	HLE7709	Call
00000003	CELQHROD	+00000256	CELQLIB	CELQHROD	HLE7709	Call
00000004	CEEOSIGG	-17F50414	CELQLIB	CEEOSIGG	HLE7709	Call
00000005	CELQHROD	+00000256	CELQLIB	CELQHROD	HLE7709	Call
00000006	main	+000000CA	*PATHNAM			Exception
00000007	CELQINIT	+00001146	CELQLIB	CELQINIT	HLE7709	Call

DSA	DSA Addr	E Addr	PU Addr	PU Offset	Comp Date	Attributes				
00000001	00000001082FABC0	00000000201C8F08	00000000201C8F08	00000000	20040312	XPLINK	EBCDIC	POSIX	Floating	Point
00000002	00000001082FD4E0	0000000020251100	0000000020251100	00000956	20040312	XPLINK	EBCDIC	POSIX	Floating	Point
00000003	00000001082FDDE0	0000000020182570	0000000020182570	00000256	20040312	XPLINK	EBCDIC	POSIX	Floating	Point
00000004	00000001082FE120	000000002024A7D8	000000002024A7D8	17F50408	20040312	XPLINK	EBCDIC	POSIX	Floating	Point
00000005	00000001082FEF40	0000000020182570	0000000020182570	00000256	20040312	XPLINK	EBCDIC	POSIX	Floating	Point
00000006	00000001082FF180	00000000200770C0	0000000000000000	*****	20040813	XPLINK	EBCDIC	POSIX	IEEE	
00000007	00000001082FF280	000000002013D010	000000002013D010	00001146	20040312	XPLINK	EBCDIC	POSIX	Floating	Point

- Note that *execution* (RMODE) is still below-the-bar
- Only downstack transitions are from operating system
- Fewer control blocks than 31-bit; SYSMDUMP/IPCS VERBX LEDATA essential.



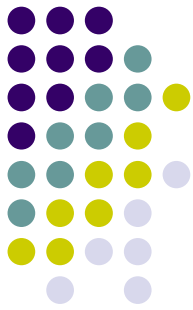
XPLink Transitions Tracing

New for z/OS V1R8

- Trace level LE=20 for the TRACE run-time option specifies that transitions between XPLink and non-XPLink be recorded (31-bit only).
 - Trace entry type 7 occurs when an XPLINK function calls a non-XPLINK function.
 - Trace entry type 8 occurs when a non-XPLINK function calls an XPLINK function.
 - Format for trace table entries 7 and 8:

ModuleNameOfCallingFunction : NameOfCallingFunction →

ModuleNameOfCalledFunction : NameOfCalledFunction

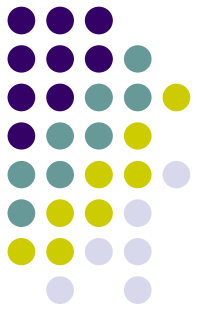


XPLink Transitions Tracing

Example:

TRACE=(ON,32K,DUMP,LE=20)

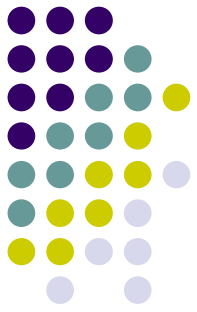
Displacement	Trace Entry in Hexadecimal	Trace Entry in EBCDIC
+000000	Time 23.04.44.393727 Date 2005.04.29 Thread ID... 2050D5F000000000	
+000010	Member ID.... 03 Flags..... 000000 Entry Type..... 00000008	
+000018	C3C5D3C8 E5F0F0F3 40404040 40404040 7AC5C4C3 E9C8C9D5 E5404040 40404040	CELHV003 :EDCZHINV
+000038	40404040 40404040 40404040 40404040 4060606E 81F8F5F9 83F0F1A7 40404040	:main -->a859c01x
+000058	40404040 7A948189 95404040 40404040 40404040 40404040 40404040 40404040	1
+000078	40404040 404040F1	
+000080	Time 23.04.44.399327 Date 2005.04.29 Thread ID... 2050D5F000000000	
+000090	Member ID.... 03 Flags..... 000000 Entry Type..... 00000007	
+000098	C3C5C5D7 D3D7D2C1 40404040 40404040 7AC3C5C5 D7C8E3D3 C3404040 40404040	CEEPLPKA :CEEPHTLC
+0000B8	40404040 40404040 40404040 40404040 4060606E C3C5C5D7 D3D7D2C1 40404040	:CEEPTLOR -->CEEPLPKA
+0000D8	40404040 7AC3C5C5 D7E3D3D6 D9404040 40404040 40404040 40404040 40404040	1
+0000F8	40404040 404040F1	
+000100	Time 23.04.44.425623 Date 2005.04.29 Thread ID... 2050D5F000000000	
+000110	Member ID.... 03 Flags..... 000000 Entry Type..... 00000007	
+000118	81F8F5F9 83F0F1A7 40404040 40404040 7A948189 95404040 40404040 40404040	a859c01x :main
+000138	40404040 40404040 40404040 40404040 4060606E 81F8F5F9 83F0F240 40404040	:func1_d1 -->a859c02
+000158	40404040 7A86A495 83F16D84 F1404040 40404040 40404040 40404040 40404040	1
+000178	40404040 404040F1	
+000180	Time 23.04.44.427092 Date 2005.04.29 Thread ID... 2050D5F000000000	
+000190	Member ID.... 03 Flags..... 000000 Entry Type..... 00000008	
+000198	81F8F5F9 83F0F240 40404040 40404040 7A86A495 83F16D84 F1404040 40404040	a859c02 :func1_d1
+0001B8	40404040 40404040 40404040 40404040 4060606E C3C5D3C8 E5F0F0F3 40404040	:printf -->CELHV003
+0001D8	40404040 7A979989 95A38640 40404040 40404040 40404040 40404040 40404040	1
+0001F8	40404040 404040F1	



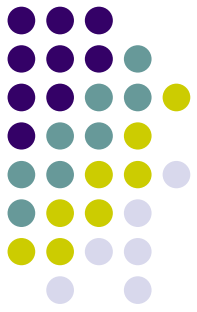
Debugging an XPLink Application

- Major points of difference between XPLink and non-XPLink
 - The Stack
 - upward-growing vs. downward-growing
 - DSA format
 - stack unwinding (backchain ptr vs. DSA size)
 - (BACKCHAIN suboption of XPLINK compile option)
 - XPLink stack ptr (GPR4) is "biased" by 0x800 bytes

Debugging an XPLink Application



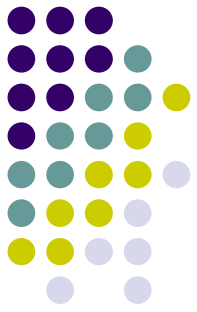
- Major points of difference between XPLink and non-XPLink continued...
 - Register conventions
 - Finding entry / return points, etc.
 - Parameter passing
 - R1 points to parm list vs. parms in regs and caller's DSA (STOREARGS suboption of XPLINK compile option)



Debugging an XPLink Application

- IPCS VERBX LEDATA
 - Similar support for tracebacks and DSAs as in CEEDUMP
- LE Storage Reporting
 - via RPTSTG Run-Time option
 - includes XPLink stack and threadstack statistics
- Full support for XPLINK-compiled functions provided by Debug Tool (**31-bit only**) and dbx debuggers
 - Through new CWIs provided by LE to traverse stack frames, locate entry points, etc.

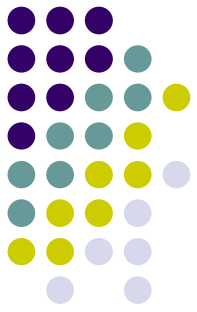
Appendix F



Callback Function Support

Reference Materials

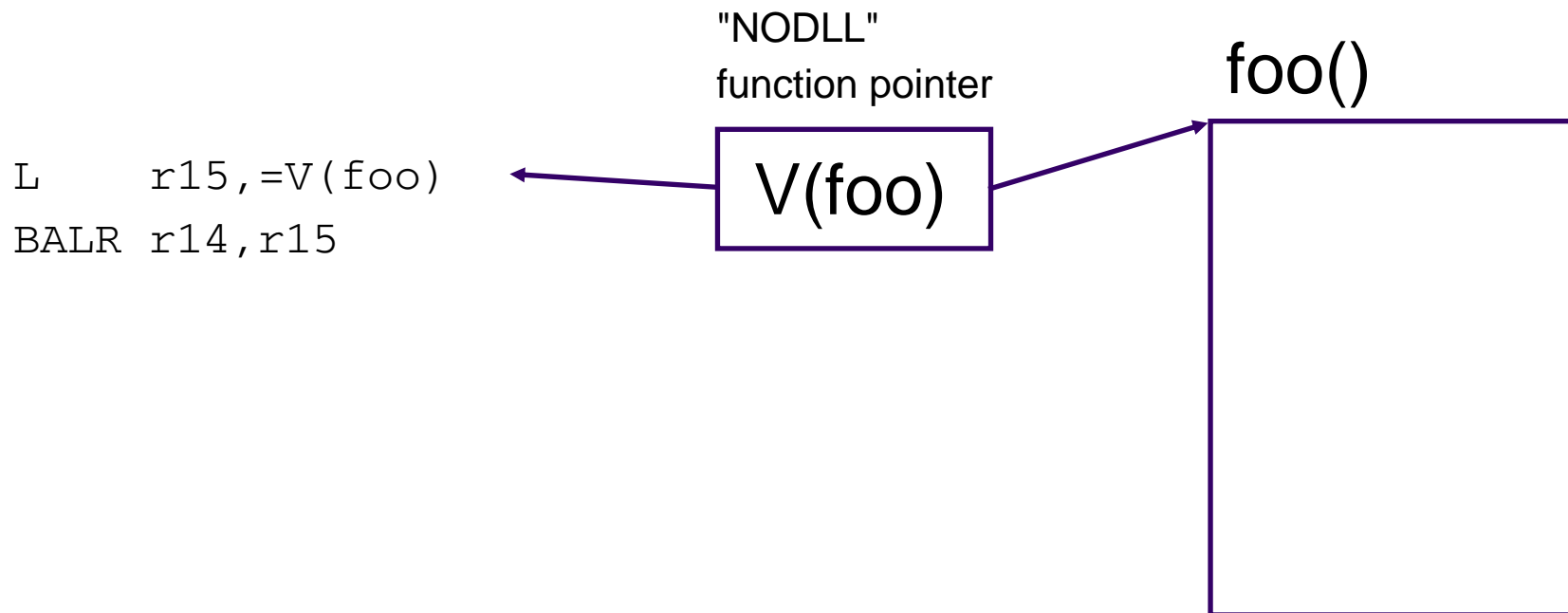


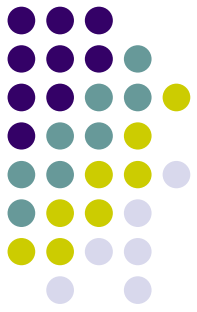


Callback Function Support

Why is `__bldxfd()` needed?

Taking the address of a function in NODLL-compiled code:

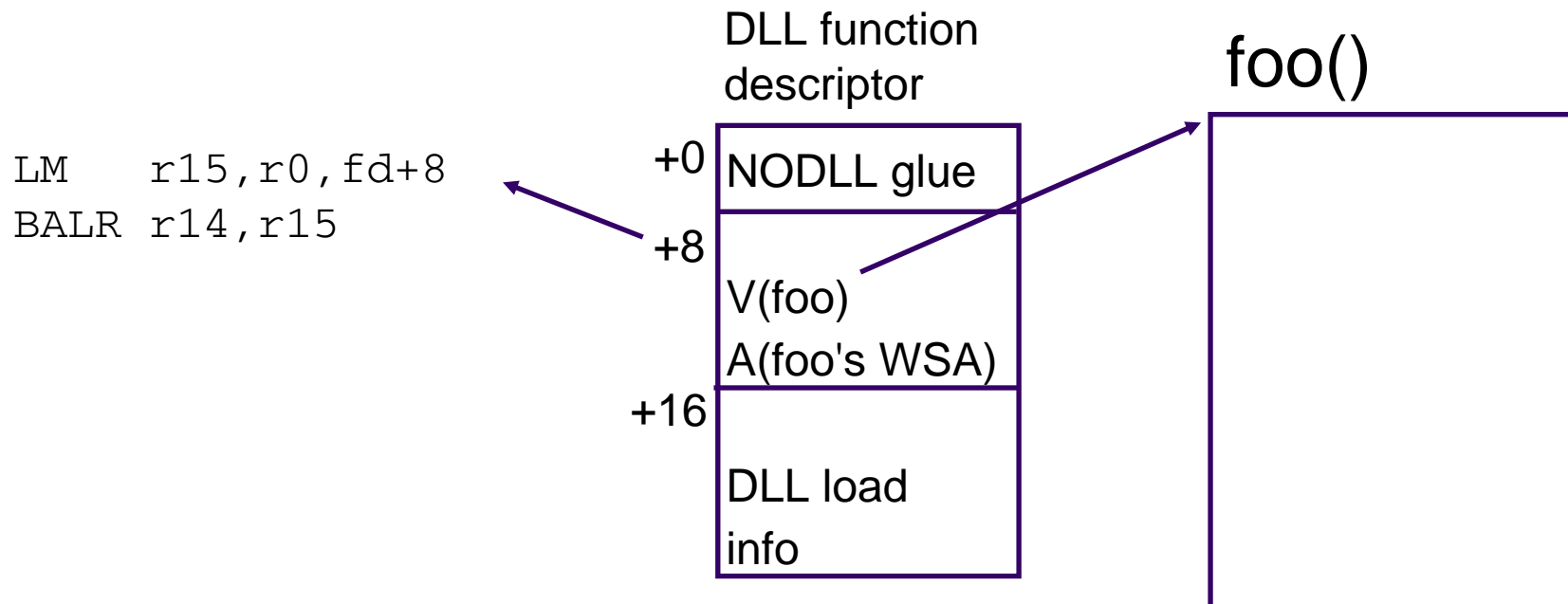


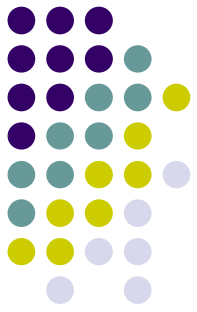


Callback Function Support...

Taking the address of a function in DLL-compiled code:

- 👉 A NODLL function pointer cannot be called from DLL-compiled code unless compiled DLL(CALLBACKANY)



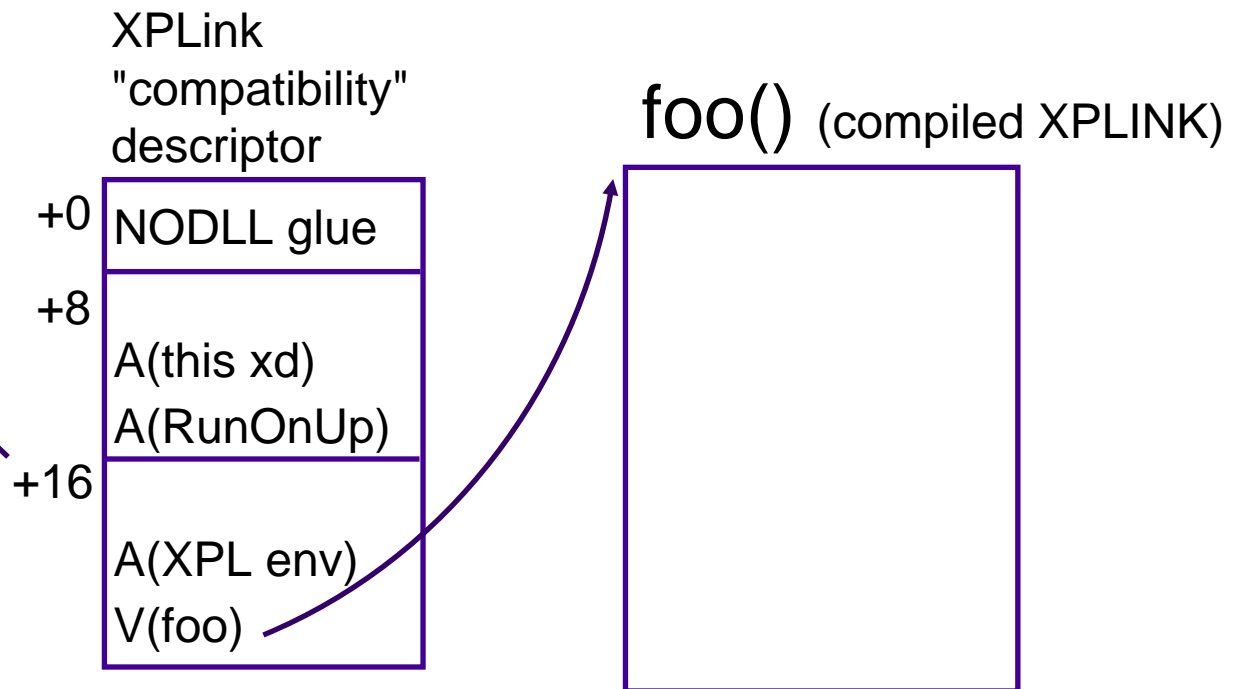


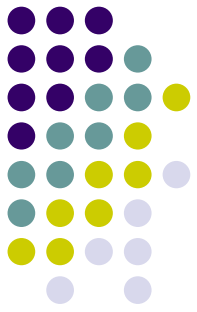
Callback Function Support...

Taking the address of a function in XPLINK-compiled code:

👉 A NODLL fp or DLL fd cannot normally be called from XPLINK-compiled code

```
LM    r5, r6, xd+16
BASR  r7, r6
NOP   ...
```



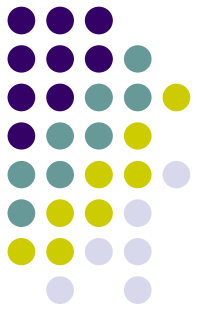


Callback Function Support...

- A NODLL function pointer or DLL function descriptor cannot normally be called from XPLINK-compiled code
- The `__bldxfd()` CWI will interrogate the input "function pointer" and convert to an XPLink compatibility descriptor if necessary
- The `__bldxfd()` CWI will be called implicitly by the compiler for each function pointer parameter passed to an exported function

```
#pragma export(foo)
typedef int (*FP)(void);
void foo(FP fpParm1, int parm2, FP fpParm3) {
```

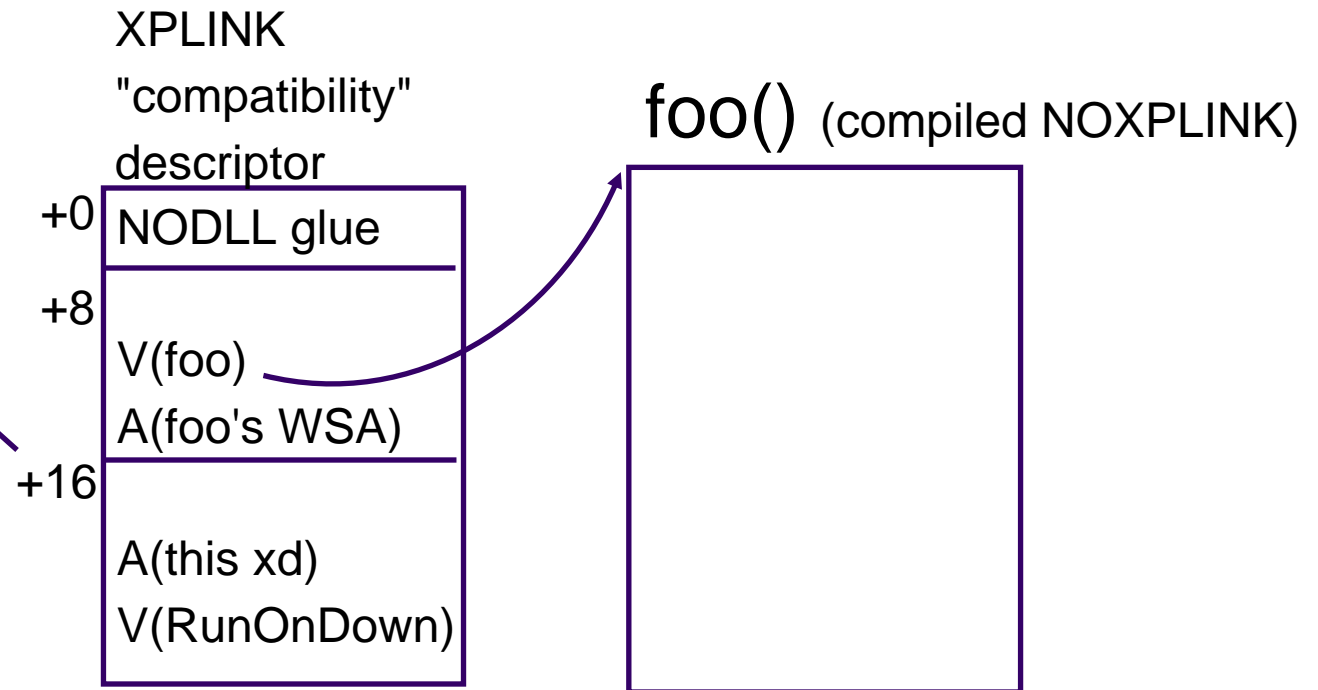
- `__bldxfd()` is *not* called for:
 - ▶ function pointers passed inside a structure, or global function pointers
- You need to call `__bldxfd()` explicitly in these cases, or see "New Compiler solutions to callback problem" in a couple of pages



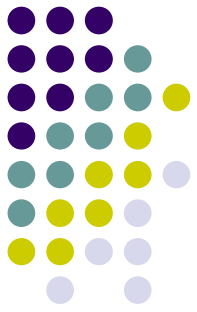
Callback Function Support...

What if we take the address of NOXPLINK-compiled code from an XPLink function? Same rules apply, but different style XPLink compatibility descriptor.

```
LM    r5, r6, xd+16  
BASR  r7, r6  
NOP   ...
```



New Compiler Solutions to Callback Problem



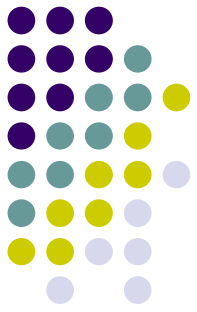
- `__callback` type cast qualifier
 - In the following example, all calls to `(*func_p)()` first result in a call to `__bldxfd()`.

```
#if !__XPLINK_CALLBACK__  
#define __callback  
#endif
```

```
...  
void (* __callback func_p) (void);  
...
```

- `XPLINK(CALLBACK)` compiler option
 - *ALL* calls through function pointers result in a call to `__bldxfd()`.
 - Not recommended for performance reasons.

64-bit solution to callback problem



Why again was `__bldxfd()` needed?

For taking the address of a function in NODLL-compiled code.

Why is `__bldxfd()` not needed in 64-bit?

Because there is no non-DLL non-XPLink code to interface with. All 64-bit function descriptors are RD-con/VD-con doubleword adcon pairs (16 bytes, loaded using LMG from zero past the descriptor location).